On the Complexity of Ordered Colourings

JAN VAN DEN HEUVEL

joint work with ARVIND GUPTA, JÁN MAŇUCH
LADISLAV STACHO & XIAOHONG ZHAO
(Simon Fraser University, Canada)

Centre for Discrete and Applicable Mathematics
Department of Mathematics
London School of Economics and Political Science

A simple ordered 3-colouring problem

- **given**: graph $G = (V, E)$
  with vertices $V = \{1, 2, \ldots, n\}$

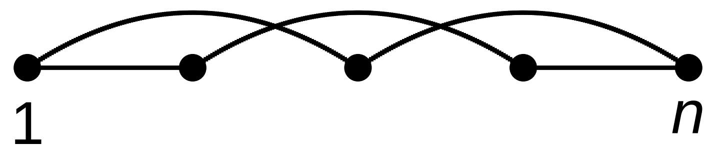
- **required**: a vertex colouring with colours $\{A, B, C\}$
  i.e., a function $\varphi : V \longrightarrow \{A, B, C\}$

  such that:

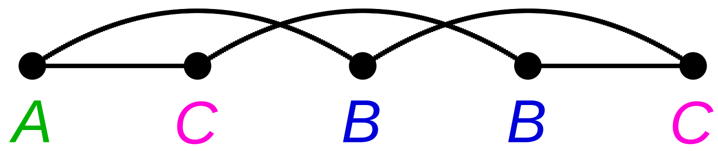
  - $\forall uv \in E : \varphi(u) \neq \varphi(v)$ (i.e., a proper colouring)
  - all vertices coloured A are smaller than all vertices coloured B

Example

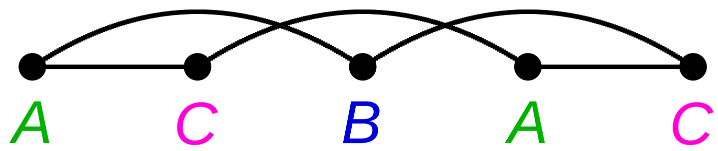
- suppose the graph is :



  - allowed colouring :



  - not allowed :



decision problem :

- given a graph G with vertices $\{1, 2, \ldots, n\}$
  how easy is it to decide if it allows such a 3-colouring ?

Deciding ordered 3-colouring

Theorem

- *deciding if a given graph allows this kind of 3-colouring can be done in **polynomial time***

sketch of proof

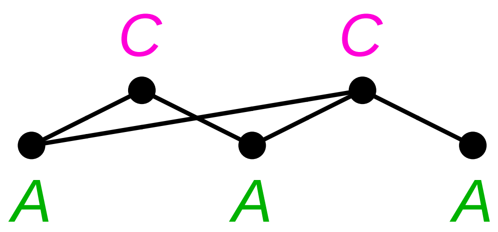
- choose a vertex m from $1, \ldots, n$ (n choices)
- assume m forms the boundary between
  the vertices that can have colour A ($\leq m$)
  and the vertices that can have colour B ($> m$)

Deciding ordered 3-colouring

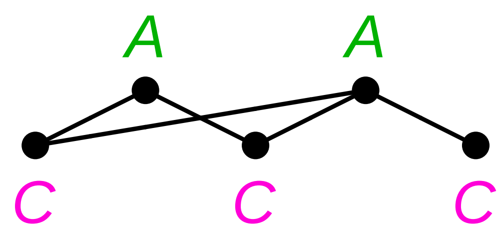
sketch of proof (cont.)

- then vertices $1, \ldots, m$ can only be coloured A or C and vertices $m+1, \ldots, n$ can only be coloured B or C
  - so check if the two parts are **bipartite**
- suppose "**yes**"
  - so $G[1, \ldots, m]$ has bipartite components

    same for $G[m+1, \ldots, n]$
- each component has 2 different choices for a 2-colouring

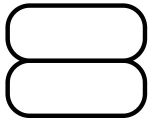
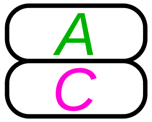


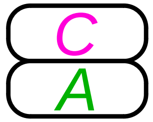
or

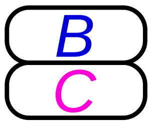


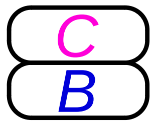
Deciding ordered 3-colouring

sketch of proof (cont.)

- for each component X_i in $G[1, \ldots, m]$
  - choose a **top** and a **bottom** part: 
  - introduce a **Boolean variable** x_i with the meaning
    - $x_i = \text{TRUE}$: 

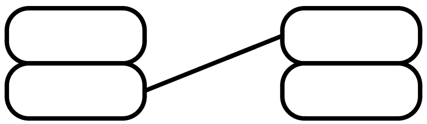
      $x_i = \text{FALSE}$: 

- for each component Y_j in $G[m+1, \ldots, n]$
  - choose a **top** and a **bottom** part
  - introduce a **Boolean variable** y_j with the meaning
    - $y_j = \text{TRUE}$: 

      $y_j = \text{FALSE}$: 

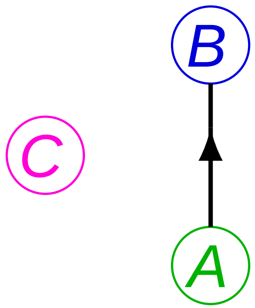
Deciding ordered 3-colouring

sketch of proof (cont.)

- x_i : $\begin{array}{|c|}\hline A \\ \hline C \\ \hline \end{array}$ $\overline{x_i}$: $\begin{array}{|c|}\hline C \\ \hline A \\ \hline \end{array}$ y_j : $\begin{array}{|c|}\hline B \\ \hline C \\ \hline \end{array}$ $\overline{y_j}$: $\begin{array}{|c|}\hline C \\ \hline B \\ \hline \end{array}$
- each edge between a component X_i and a component Y_j gives rise to a **2-literal disjunction**
  - **e.g.** :  $\longleftrightarrow$ $\overline{x_i} \vee y_j$
- doing this for **all edges** between X_i and Y_j components gives a **conjunction of 2-literal disjunctions**
- i.e., an instance of the **2-SAT** problem
  - well-known to be **polynomial time** decidable

Generalising the problem

- a way to look at this colouring problem is as if there is a **order relation** on the colours $\{A, B, C\}$

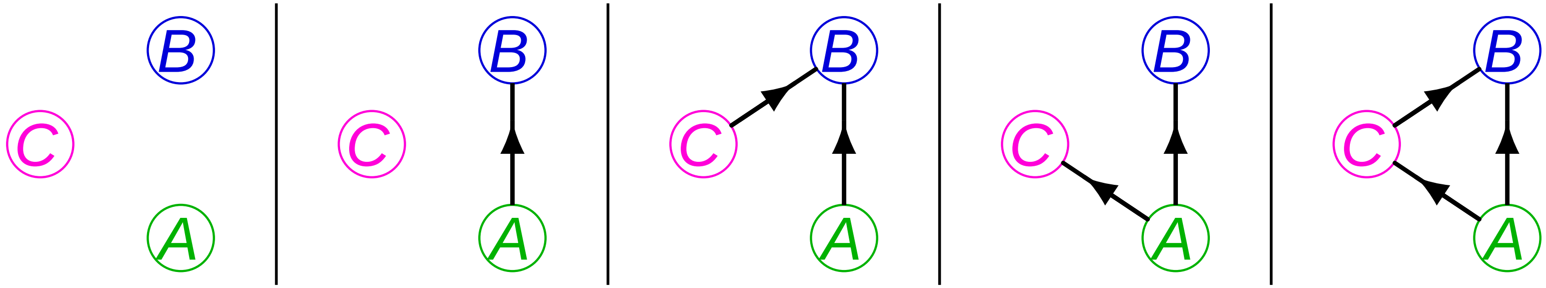


means : "all A before all B"
"A - C and B - C are indifferent"

- what would happen if we consider other order relations on the colours ?

More general ordered 3-colourings

- possible posets with 3 elements :



Theorem

- ***deciding*** *if a certain graph with vertices* $1, \ldots, n$ *has a colouring according to these posets is*
  - ***NP-complete*** *for the first poset (just 3-colouring)*
  - ***polynomial*** *for the others*

General ordered colourings with more colours

- Rule 1 : $A \longrightarrow B$ means : "all A before all B"
- suppose we allow **any number** of colours and **any poset** on the set of colours

Theorem

- *deciding if a given graph with vertices $1, \ldots, n$ has a colouring according to a **fixed colour poset** is*
  - ***NP-complete***
    if the poset has an anti-chain of length at least 3
  - ***polynomial** otherwise*

A different ordered 3-colouring problem

- **given**: graph $G = (V, E)$ with n vertices
  vertices are numbered from $1, \ldots, n$

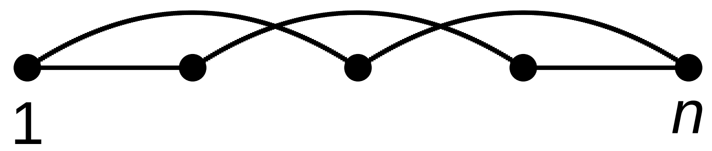
- **required**: a vertex colouring with colours $\{A, B, C\}$
  i.e., a function $\varphi : V \to \{A, B, C\}$

such that:

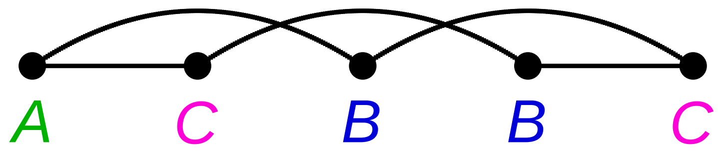
- $\forall uv \in E \ : \varphi(u) \neq \varphi(v)$ (i.e., a proper colouring)
- vertices coloured A are smaller than vertices coloured B **only for adjacent vertices**

Example

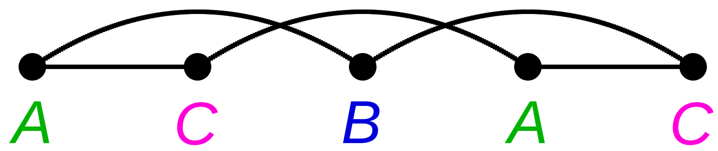
- suppose the graph is :



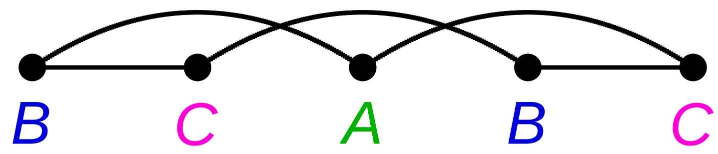
  - allowed colouring :



  - also allowed :



  - not allowed :



A different ordered 3-colouring problem

Theorem

- *this new 3-colouring problem is* ***NP-complete***

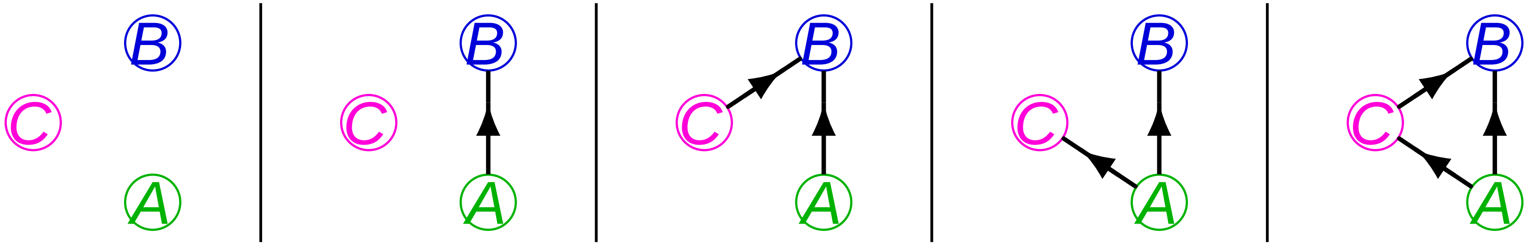
ideas of proof

- constraint on edges only means “local constraints”
  - allows to construct gadgets
- reduction to SAT

Generalising the 2nd variant for 3 colours

- <u>Rule 2</u>: $A \longrightarrow B$ means: "*A* before *B* **on edges**"

- possible posets with 3 elements:



<u>Theorem</u>

- ***deciding*** *if a certain graph with vertices* $1, \ldots, n$ *has a colouring according to Rule 2 is*
  - ***NP-complete*** *for the first two posets*
  - ***polynomial*** *for the others*

Generalising the 2nd variant with more colours

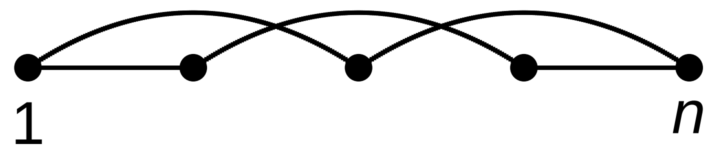
- Rule 2: $(A) \longrightarrow (B)$ means: "A before B **on edges**"
- suppose we allow **any number** of colours and **any poset** on the set of colours

Theorem

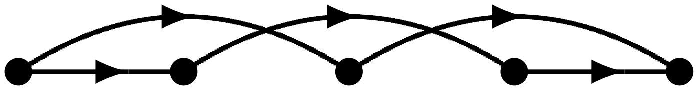
- *deciding if a given graph with vertices $1, \ldots, n$ has a colouring according to a* ***fixed colour poset*** *and Rule 2 is*
  - ***polynomial*** *if there is* ***at most one pair of non-comparable colours***
  - ***NP-complete*** *otherwise*

From colourings to homomorphisms

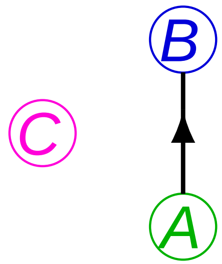
- given graph G



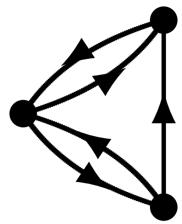
form $\overrightarrow{G}$:



- given colour poset



form $\overrightarrow{H}$:



- **then** :

G is colourable with the given poset using Rule 2

$$\Longleftrightarrow$$

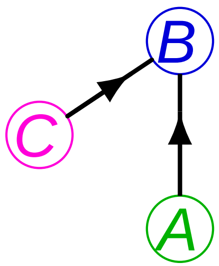
*there is a **homomorphism** $\overrightarrow{G} \longrightarrow \overrightarrow{H}$*

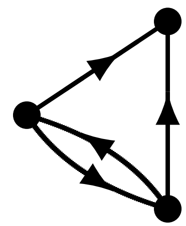
Some technical definitions

- semi-complete digraph :
  between any vertex pair there is one or two (opposite) arcs

- quasi-acyclic digraph :
  removal of all 2-cycles gives an acyclic digraph

- starting from a colour poset,
  always gives a quasi-acyclic semi-complete digraph $\vec{H}$

- e.g.: 

  gives $\vec{H}$: 

Complexity of directed homomorphism

- $\overrightarrow{H}$-COLOURING

Input: directed graph $\overrightarrow{G}$

Question: is there a homomorphism $\overrightarrow{G} \longrightarrow \overrightarrow{H}$?

Theorem (Bang-Jensen, Hell & MacGillivray, 1988)

- *let $\overrightarrow{H}$ be a semi-complete digraph then the complexity of $\overrightarrow{H}$-COLOURING is*
  - ***polynomial** if $\overrightarrow{H}$ contains **at most one directed cycle***
  - ***NP-complete** otherwise*

Complexity of acyclic directed homomorphism

- ACYCLIC-$\overrightarrow{H}$-COLOURING

  Input : acyclic directed graph $\overrightarrow{G}$

  Question : is there a homomorphism $\overrightarrow{G} \longrightarrow \overrightarrow{H}$?

Theorem

- *let $\overrightarrow{H}$ be a quasi-acyclic semi-complete digraph then the complexity of ACYCLIC-$\overrightarrow{H}$-COLOURING is*
  - ***polynomial** if $\overrightarrow{H}$ contains **at most one 2-cycle***
  - ***NP-complete** otherwise*

Open problems

- is there a characterisation for semi-complete digraphs $\overrightarrow{H}$ on the complexity of ACYCLIC-$\overrightarrow{H}$-COLOURING ?

  (probably “yes”)

- what about a characterisation for general digraphs $\overrightarrow{H}$ on the complexity of ACYCLIC-$\overrightarrow{H}$-COLOURING ?

  (likely to be very hard; open even if we restrict $\overrightarrow{H}$ to trees)

- the Holy Grail :
  characterisation for general digraphs $\overrightarrow{H}$ on the complexity of $\overrightarrow{H}$-COLOURING