

# Distance-Two Colouring of Graphs

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joint work with

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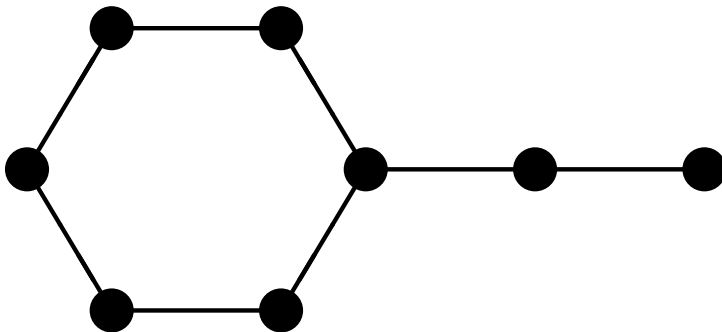


## *The start*

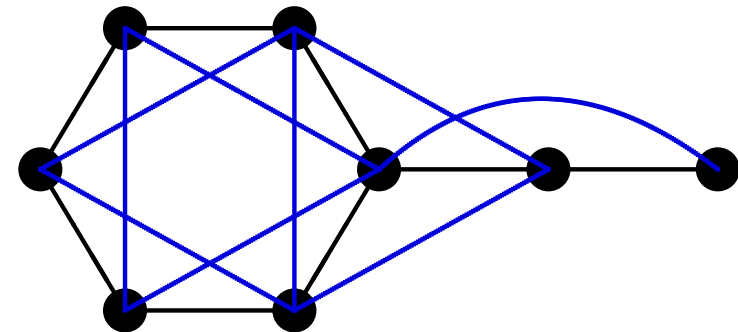
- on way to look at **vertex-colouring** :  
vertices at **distance one** must receive different colours
- suppose we require vertices at **larger distances**  
to receive different colours as well
- for today we only look at **distance two**

# The square of a graph

- distance-two colouring can be modelled using the **square  $G^2$  of a graph**:
  - same vertex set as  $G$
  - edges between vertices with **distance at most 2 in  $G$**   
( = **are adjacent** or **have a common neighbour** )



$G$



$G^2$

# Colouring the square of a graph

## Easy facts

- $\chi(G^2) \geq \Delta(G) + 1$

and

- $\Delta(G^2) \leq \Delta(G)^2$ , so  $\chi(G^2) \leq \Delta(G)^2 + 1$

## Question

- is the **upper bound** relevant?

- there are **at most 4 graphs** with  $\chi(G^2) = \Delta(G)^2 + 1$

- and **infinitely many graphs** with

$$\chi(G^2) = \Delta(G)^2 - \Delta(G) + 1$$

# The square of planar graphs

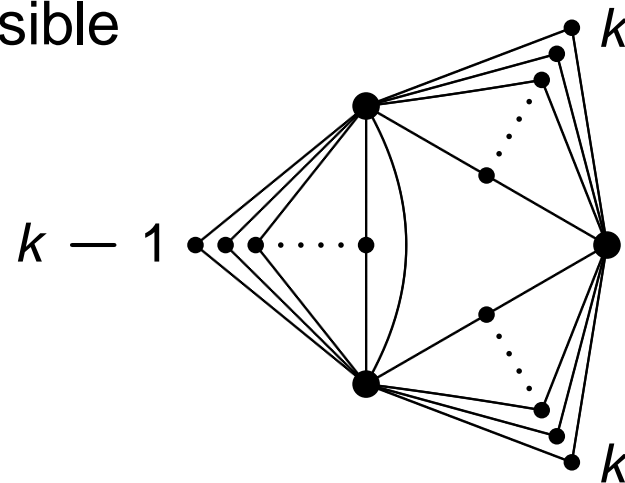
Conjecture (Wegner, 1977)

■  $G$  planar

$$\implies \chi(G^2) \leq \begin{cases} 7, & \text{if } \Delta = 3 \\ \Delta + 5, & \text{if } 4 \leq \Delta \leq 7 \\ \lfloor \frac{3}{2} \Delta \rfloor + 1, & \text{if } \Delta \geq 8 \end{cases}$$

■ bounds would be best possible

case  $\Delta = 2k \geq 8$ :



## What is known for large $\Delta$

$G$  planar  $\implies$

■  $\chi(G^2) \leq 8\Delta - 22$  (Jonas, PhD, 1993)

■  $\chi(G^2) \leq 3\Delta + 5$  (Wong, MSc, 1996)

■  $\chi(G^2) \leq 2\Delta + 25$  (vdH & McGuinness, 2003)

■  $\chi(G^2) \leq \lceil \frac{9}{5}\Delta \rceil + 1$  (for  $\Delta \geq 47$ )  
(Borodin, Broersma, Glebov & vdH, 2001)

■  $\chi(G^2) \leq \lceil \frac{5}{3}\Delta \rceil + 24$  (for  $\Delta \geq 241$ )  
(Molloy & Salavatipour, 2005)

## First new results

### Theorem

$$\blacksquare \quad G \text{ planar} \implies \chi(G^2) \leq \left(\frac{3}{2} + o(1)\right) \Delta \quad (\Delta \rightarrow \infty)$$

we actually prove the **list-colouring version** for much larger classes of graphs :

### Theorem

$$\blacksquare \quad \text{graph } G \text{ } K_{3,k}\text{-minor free for some fixed } k \implies ch(G^2) \leq \left(\frac{3}{2} + o(1)\right) \Delta$$

## *Even more general results ?*

### Property

- graph  $G$   $H$ -minor free for some fixed graph  $H$   
 $\implies ch(G^2) \leq C_H \Delta$  for some constant  $C_H$

### Question

- given  $H$ , what is the best  $C_H$  for large  $\Delta$ ?

e.g.

- for  $H = K_5$  we have  $2 \leq C_{K_5} \leq 9$



# The clique number

## Corollary

- graph  $G$   $K_{3,k}$ -minor free for some fixed  $k$   
 $\implies \omega(G^2) \leq \left(\frac{3}{2} + o(1)\right) \Delta$

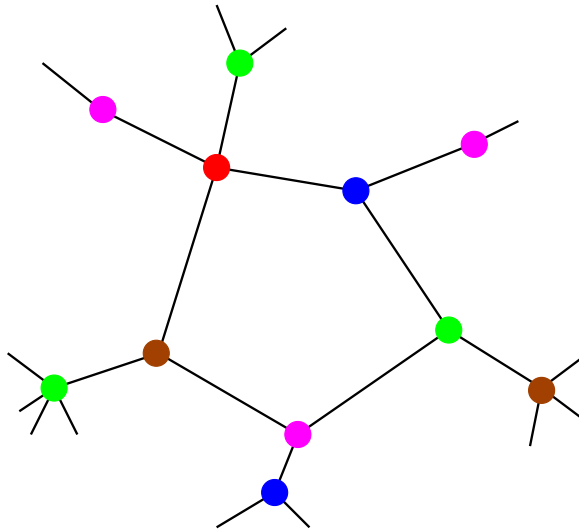
this can be partially improved to

## Theorem

- $G$  planar  $\implies \omega(G^2) \leq \frac{3}{2} \Delta + O(1)$

## *A related (?) problem*

- **plane graph**: planar graph with a given embedding
- **cyclic colouring of a plane graph**:
  - vertex-colouring so that
  - vertices incident to the same **face** get a different colour



## *A related (?) problem*

- **plane graph**: planar graph with a given embedding
- **cyclic colouring of a plane graph**:
  - vertex-colouring so that
  - vertices incident to the same **face** get a different colour
- **cyclic chromatic number**  $\chi^*(G)$ :  
minimum number of colours needed for a cyclic colouring
- $\Delta^*(G)$ : size of largest face of  $G$

### Easy

- $G$  plane graph  $\implies \chi^*(G) \geq \Delta^*(G)$

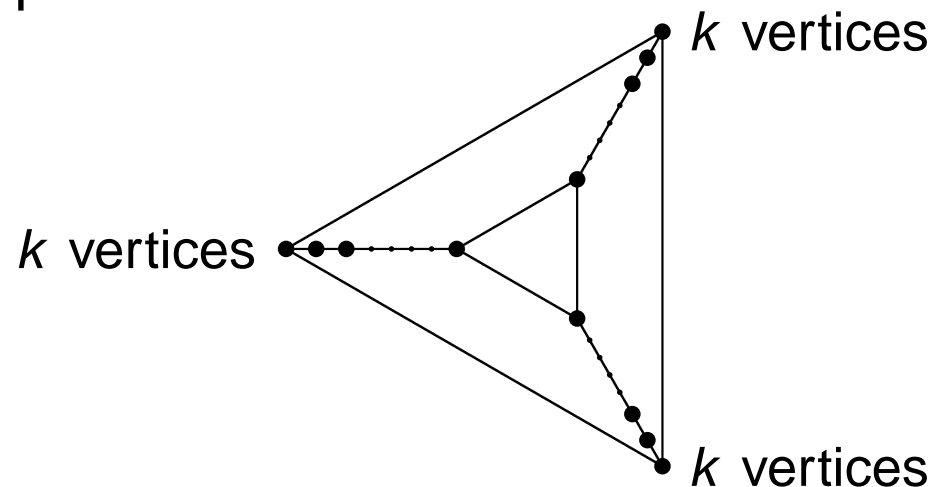
## Conjecture for the related(?) problem

Conjecture (Borodin, 1984)

■  $G$  plane graph  $\implies \chi^*(G) \leq \frac{3}{2} \Delta^*(G)$

■ bound would be best possible

case  $\Delta^* = 2k$ :



# Bounds on the cyclic chromatic number

$G$  plane graph  $\implies$

■  $\chi^*(G) \leq 2 \Delta^*(G)$  (Ore & Plummer, 1969)

■  $\chi^*(G) \leq \frac{9}{5} \Delta^*(G)$  (Borodin, Sanders & Zhao, 1999)

■  $\chi^*(G) \leq \frac{5}{3} \Delta^*(G)$  (Sanders & Zhao, 2001)

## Theorem

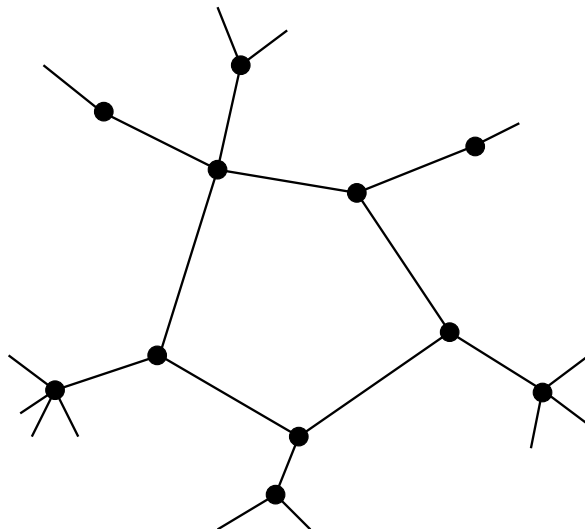
■  $G$  plane graph

$$\implies \chi^*(G) \leq \left(\frac{3}{2} + o(1)\right) \Delta^* \quad (\Delta^* \rightarrow \infty)$$

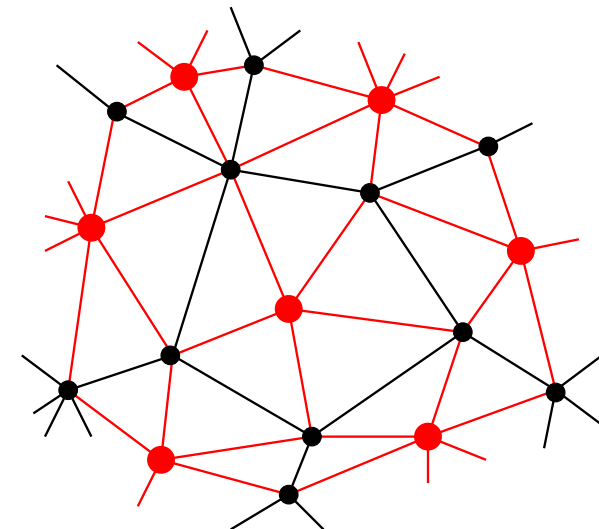
(in fact, we again prove the **list-colouring version**)

# From cyclic colouring to “distance-two” colouring

- given a plane graph  $G$ , form new graph  $G^F$  by
  - adding a vertex in each face
  - adding edges from new vertex to all vertices of the face

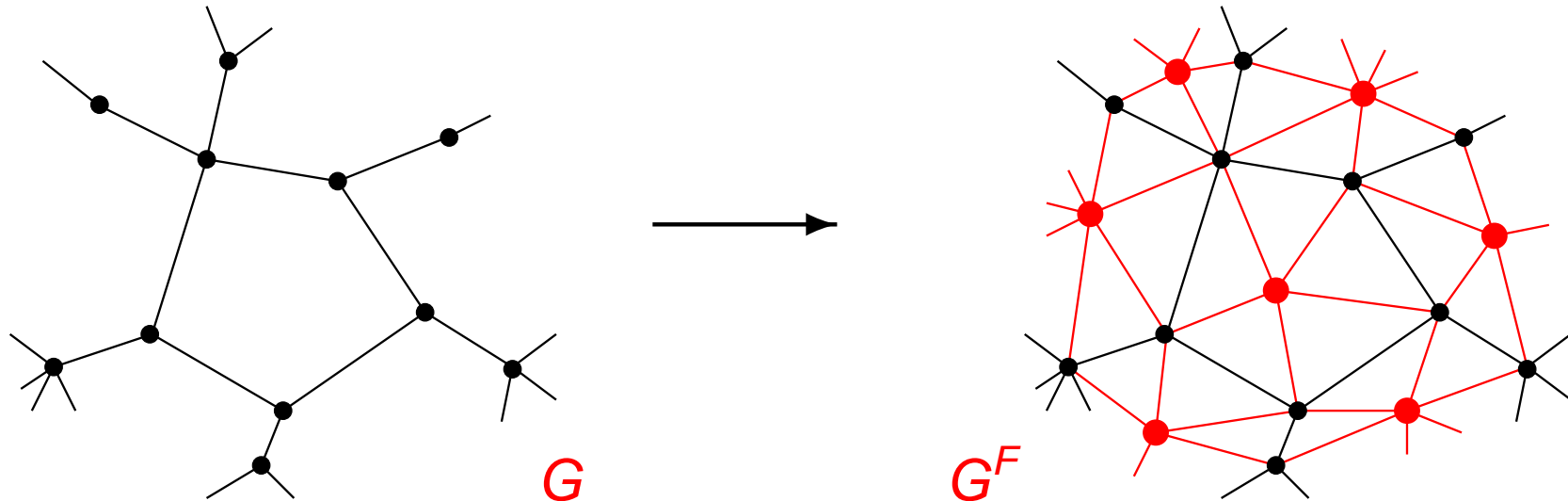


$G$



$G^F$

# From cyclic colouring to “square” colouring



- colouring the square of  $G^F$  gives cyclic colouring of  $G$ 
  - but also colours the faces ( not asked for, but o.k. )
  - and degrees of old vertices may be more than  $\Delta^*(G)$   
( serious problem )

## ***Both in one go***

**idea** : do not treat all vertices equal :

- some vertices **need to be coloured**
- some vertices **determine distance two**
  
- vertices can be of both types
  - as are all vertices when colouring the square
- or the types can be disjoint
  - as for “faces” and “vertices” for cyclic colouring
- or anything in between



## Formalising our clever little idea

- **given**: graph  $G$ , subsets  $A, B \subseteq V(G)$
- **$(A, B)$ -colouring of  $G$** : colouring of vertices in  $B$  so that
  - adjacent vertices get different colours
  - vertices with a common neighbour in  $A$   
get different colours
- **list  $(A, B)$ -colouring of  $G$** :
  - similar, but each vertex in  $B$  has its own list
- $\chi(G; A, B)$  /  $ch(G; A, B)$ :  
minimum number of colours needed /  
minimum size of each list needed

## The first baby steps

- $\chi(G) = \chi(G; \emptyset, V)$
- $\chi(G^2) = \chi(G; V, V)$
- $\chi^*(G) = \chi(G^F; \text{“faces”}, \text{“vertices”})$
  
- relevant “**degree**”:  $d_B(v) =$  number of neighbours in  $B$
- “**maximum degree**”:  
 $\Delta(G; A, B) =$  maximum of  $d_B(v)$  over all  $v \in A$

### Easy

- $\Delta(G; A, B) \leq \chi(G; A, B)$   
 $\leq \Delta(G) + \Delta(G; A, B) \cdot (\Delta(G; A, B) - 1) + 1$

# The obvious(?) conjecture & results

## Conjecture

- $G$  planar,  $A, B \subseteq V(G)$ ,  $\Delta(G; A, B)$  large enough  
 $\implies \chi(G; A, B) \leq \frac{3}{2} \Delta(G; A, B) + 1$

## Theorem

- $G$  planar,  $A, B \subseteq V(G)$   
 $\implies ch(G; A, B) \leq \left(\frac{3}{2} + o(1)\right) \Delta(G; A, B)$

## Corollary

(asymptotic list version of Wegner's and Borodin's Conjecture)

- $G$  planar  $\implies ch(G^2) \leq \left(\frac{3}{2} + o(1)\right) \Delta(G)$
- $G$  plane graph  $\implies ch^*(G) \leq \left(\frac{3}{2} + o(1)\right) \Delta^*(G)$

# Sketch of the proof of square of planar graph

uses induction on the number of vertices

- **2-neighbour**: vertex at distance one or two
- $d^2(v)$ : number of 2-neighbours of  $v$   
= number of neighbours of  $v$  in  $G^2$
- we would like to remove a vertex  $v$  with  $d^2(v) \leq \frac{3}{2} \Delta$ 
  - but that can change distances in  $G - v$
- contraction to a neighbour  $u$  will solve the distance problem
  - but may increase maximum degree if  $d(u) + d(v) > \Delta$
- easy induction possible if there is an edge  $uv$   
with  $d(u) + d(v) \leq \Delta$  and  $d^2(v) \leq \frac{3}{2} \Delta$

## When easy induction is not possible

**S**, **small** vertices: degree at most **some constant  $C$**

**B**, **big** vertices: degree more than  $C$

**H**, **huge** vertices: degree at least  $\frac{1}{2} \Delta$

- small vertices have a least two big neighbours

(otherwise for those  $v$ :  $d^2(v) \leq \frac{3}{2} \Delta$ )

- a planar graph has fewer than  $3|V|$  edges  
and fewer than  $2|V|$  edges if it is bipartite

**so:**

- all but  $O(|V|/C)$  vertices are small

- fewer than  $2|B|$  vertices in  $V \setminus B$

have more than two neighbours in  $B$

## *When easy induction is not possible*

**S**, **small** vertices : degree at most **some constant C**

**B**, **big** vertices : degree more than **C**

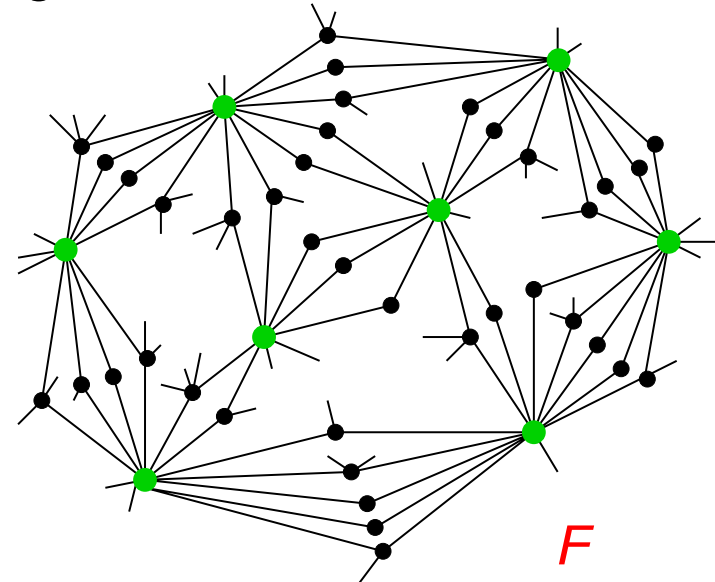
**H**, **huge** vertices : degree at least  $\frac{1}{2} \Delta$

**so :**

- “most” vertices are small
- and these have **exactly two big neighbours**  
( in fact two **huge neighbours** )

## The structure so far

- there is a **subgraph  $F$**  of  $G$  looking like :



- **green vertices  $X$**   
have degree at least  $\frac{1}{2} \Delta$
- **black vertices  $Y$**   
have degree at most  $C$
- all other neighbours of  $Y$ -vertices are also small

we can guarantee additionally :

- only “few” edges from  $X$  to rest of  $G$
- $F$  satisfies “some edge density condition”

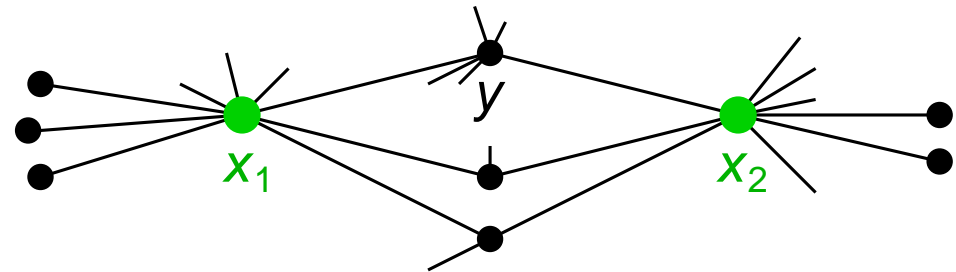
## *The other induction step*

- remove the vertices from  $Y$  (using contraction)
- colour the smaller graph (which is possible by induction)



## The other induction step

- remove the vertices from  $Y$  (using contraction)
- colour the smaller graph (which is possible by induction)
- what to do with the uncoloured  $Y$ -vertices?



- $y$  has
  - a “lot” of 2-neighbours in  $Y$  via  $x_1, x_2$
  - at most  $C^2$  other 2-neighbours in  $Y$
  - at most  $(d_G(x_1) - d_F(x_1)) + (d_G(x_2) - d_F(x_2))$  2-neighbours outside  $Y$  via  $x_1, x_2$
  - at most  $C^2$  other 2-neighbours outside  $Y$

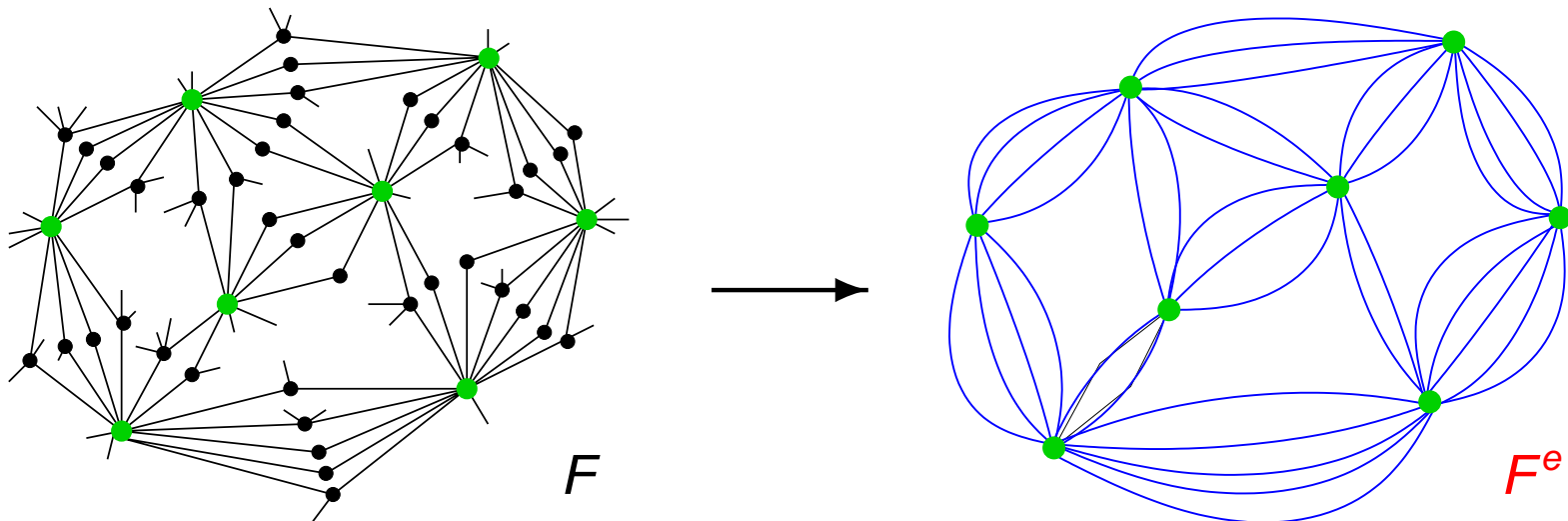
## Transferring to edge-colouring

- so a vertex  $y$  from  $Y$  has at least

$$\left(\frac{3}{2} + \varepsilon\right) \Delta - (d_G(x_1) - d_F(x_1)) - (d_G(x_2) - d_F(x_2)) - C^2$$

colours still available

- colouring  $Y$  is “almost” like list-colouring edges of the multigraph  $F^e$  :



## Edge-colouring multigraphs

- $\chi'(G)$  : chromatic index of multigraph  $G$
- $ch'(G)$  : list chromatic index of multigraph  $G$
- $\chi'_F(G)$  : fractional chromatic index of multigraph  $G$

Theorem (Kahn, 1996, 2000)

- $G$  multigraph, with  $\Delta$  large enough  
 $\implies ch'(G) \approx \chi'(G) \approx \chi'_F(G)$
- in fact, Kahn's proofs provide something much more general

# Kahn's result

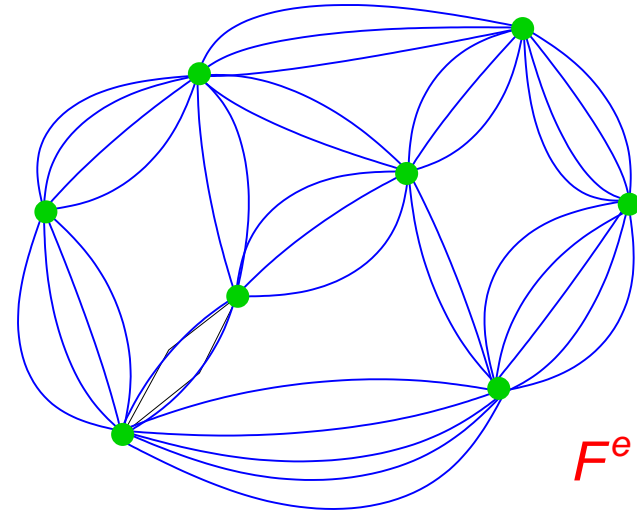
## Theorem (Kahn, 2000)

for  $0 < \delta < 1$ ,  $\alpha > 0$  there exists  $\Delta_{\delta, \alpha}$  so that if  $\Delta \geq \Delta_{\delta, \alpha}$ :

- $G$  a multigraph with maximum degree  $\Delta$
- each edge  $e$  has a list  $L(e)$  of colours so that
  - for all edges  $e$ :  $|L(e)| \geq \alpha \Delta$
  - for all vertices  $v$ : 
$$\sum_{e \ni v} |L(e)|^{-1} \leq 1 \cdot (1 - \delta)$$
  - for all  $K \subseteq G$  with  $|V(K)| \geq 3$  odd:
$$\sum_{e \in E(K)} |L(e)|^{-1} \leq \frac{1}{2} (|V(K)| - 1) \cdot (1 - \delta)$$
- then there exists a proper colouring of the edges of  $G$   
so that each edge gets colours from its own list

## Kahn's approach for our case

- we have a multigraph  $F^e$  :

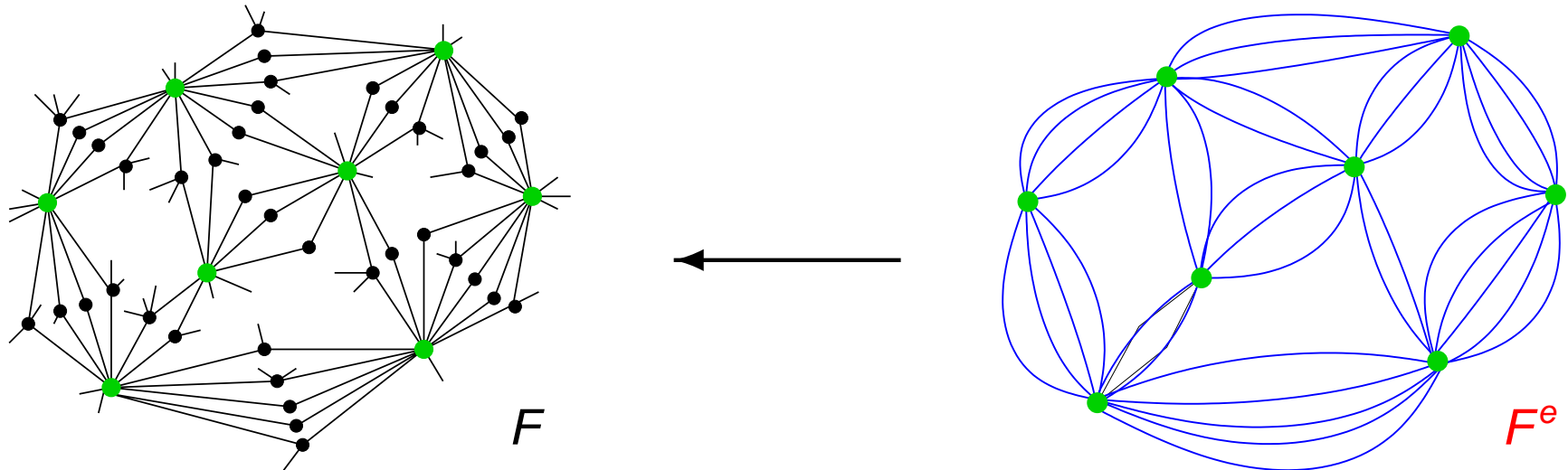


- so that each edge  $e = x_1x_2$  has a list  $L(e)$  of at least  
 $(\frac{3}{2} + \varepsilon) \Delta - (d_G(x_1) - d_F(x_1)) - (d_G(x_2) - d_F(x_2)) - C^2$   
colours
- and  $F^e$  satisfies “some edge density condition”

## Extending Kahn's approach

- these conditions guarantee that Kahn's conditions are satisfied for  $F^e$

- $\implies$  we can edge-colour  $F^e$



- $\implies$  we can colour the  $Y$ -vertices in  $F$   
choosing from the left-over colours for each

- also: we can deal with the “almost” list-edge colouring

## Some open problems

- prove the next step :

$$G \text{ planar} \implies \chi(G^2) \leq \frac{3}{2} \Delta + O(1)$$

- for  $G$  planar,  $\Delta \leq 3$  we know :

- $\chi(G^2) \leq 7$  ( Thomassen, 2007 )

- $ch(G^2) \leq 8$  ( Cranston & Kim, 2006 )

- what is the right upper bound for  $ch(G^2)$  in this case ?

- Wegner's Conjecture for  $4 \leq \Delta \leq 7$  :

$$G \text{ planar} \implies \chi(G^2) \leq \Delta + 5 ?$$

# Distance-two colouring and edge-colouring

- is there some deep relation between edge-colouring multigraphs and distance-two colouring graphs ?
- our proof uses edge-colouring to answer questions on distance-two colouring

- Theorem (Shannon, 1949)

for any multigraph  $G \implies \chi'(G) \leq \frac{3}{2} \Delta$

- Conjecture (Wegner, 1977)

for a planar graph  $G$ ,  $\Delta \geq 8 \implies \chi(G^2) \leq \frac{3}{2} \Delta + 1$



# Edge-colouring and distance-two colouring

Conjecture 1 (Vizing, 1975)

- for any multigraph  $G \implies ch'(G) = \chi'(G)$

Conjecture 2 (Kostochka & Woodall, 2001)

- for any graph  $G \implies ch(G^2) = \chi(G^2)$
- Conjecture 1 is known for **bipartite graphs** (Galvin, 1995)  
proof doesn't require knowledge of  $\chi'(G)$  and  $ch'(G)$
- Conjecture 2 is only known for some special graph classes,  
usually since we know **both  $\chi(G^2)$  and  $ch(G^2)$**