

Graph Colouring with Distances

JAN VAN DEN HEUVEL

Department of Mathematics
London School of Economics and Political Science

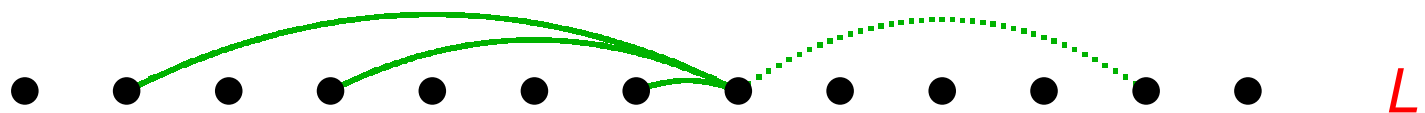


The basics of graph colouring

- **vertex-colouring** with k colours:
adjacent vertices must receive different colours
- **chromatic number** $\chi(G)$:
minimum k such that a vertex-colouring exists

Some essential graph parameters

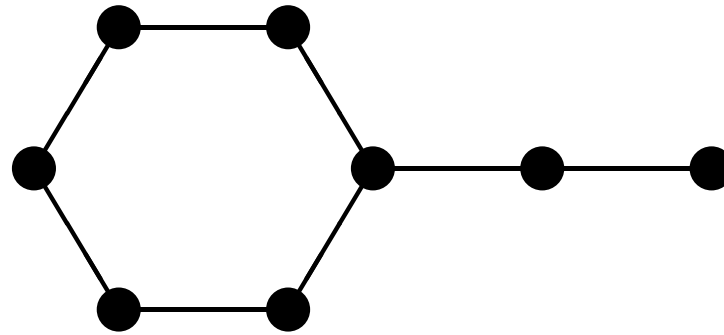
- $\delta(G)$: minimum vertex degree
- $\Delta(G)$: maximum vertex degree
- ■ G is **k -degenerate**: every subgraph of G has minimum degree at most k
- equivalent:
there is an ordering L of the vertices of G , such that every vertex has at most k neighbours that come earlier in the ordering



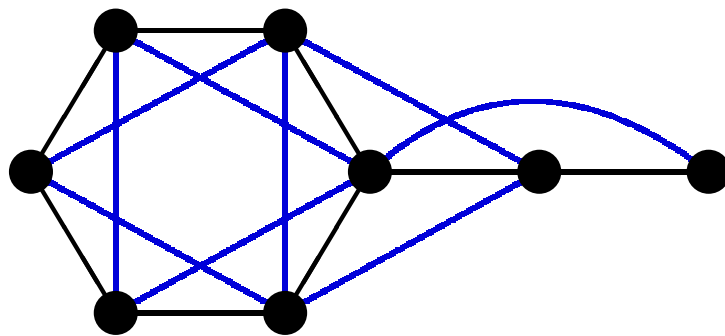
Another way to look at vertex-colouring

- vertex-colouring:
 - vertices at distance one must receive different colours
- now suppose we want vertices at larger distances (say, up to distance d) to receive different colours as well
- can be modelled using the d -th power G^d of a graph:
 - same vertex set as G
 - edges between vertices with distance at most d in G

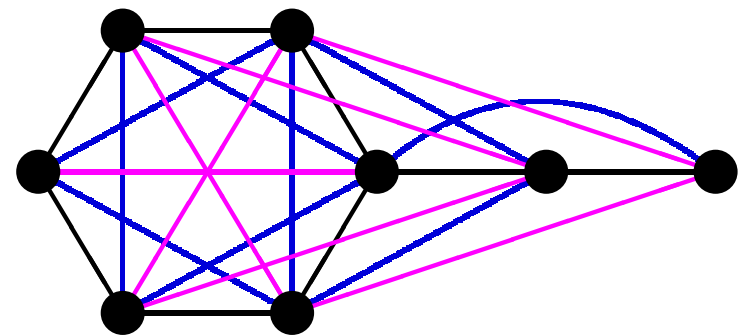
Powers of a graph



G

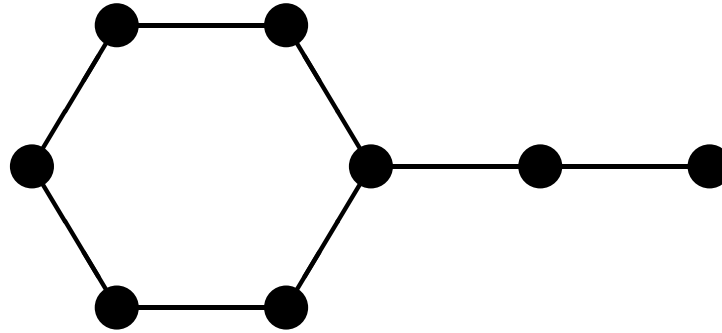


G^2

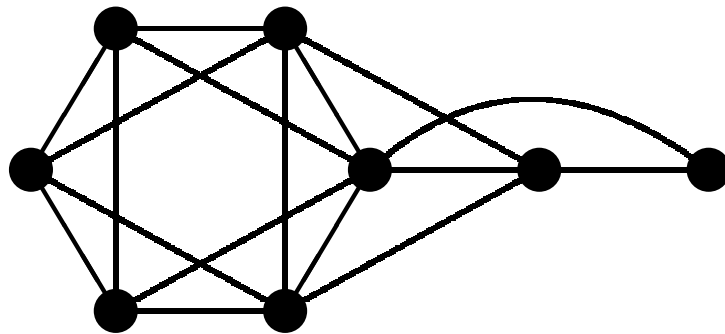


G^3

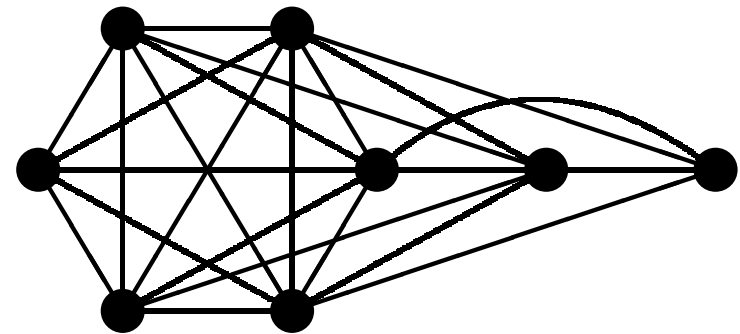
Powers of a graph



G



G^2



G^3

Colouring powers of a graph

easy facts

■ $d \geq 2 \implies \chi(G^d) \geq \Delta(G) + 1$

■ and $\chi(G^d) \leq 1 + \sum_{i=0}^{d-1} \Delta(G) (\Delta(G) - 1)^i$

■ for **connected** graphs, we have **equality** of the **upper bound** only if

■ any d : **odd cycles** C_{2d+1}

■ $d = 2$: C_5 and **two or three** more graphs
(including **Petersen graph**)

The square of k -degenerate graphs

fairly easy

■ G k -degenerate

$\implies G^2$ is $((2k - 1) \Delta(G))$ -degenerate

so

■ G planar $\implies \chi(G^2) \leq 9 \Delta(G) + 1$

The square of planar graphs

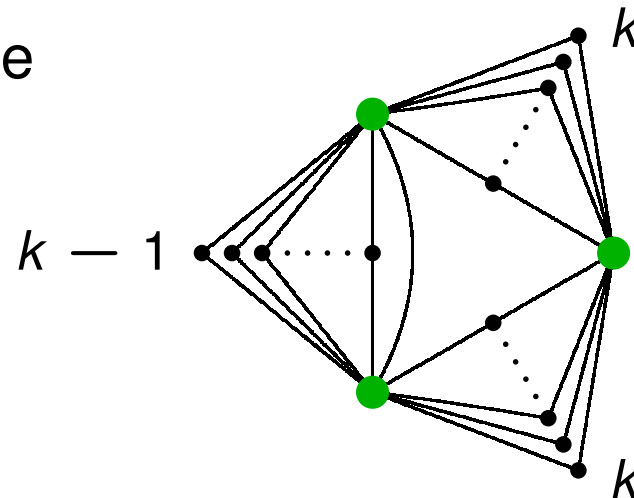
Conjecture (Wegner, 1977)

■ G planar

$$\implies \chi(G^2) \leq \begin{cases} 7, & \text{if } \Delta(G) = 3 \\ \Delta(G) + 5, & \text{if } 4 \leq \Delta(G) \leq 7 \\ \lfloor 3/2 \Delta(G) \rfloor + 1, & \text{if } \Delta(G) \geq 8 \end{cases}$$

■ bounds would be best possible

case $\Delta(G) = 2k \geq 8$:



Towards Wegner's Conjecture

G planar \implies

■ $\chi(G^2) \leq 8 \Delta(G) - 22$ (Jonas, PhD, 1993)

■ $\chi(G^2) \leq 3 \Delta(G) + 5$ (Wong, MSc, 1996)

■ $\chi(G^2) \leq 2 \Delta(G) + 25$ (vdH & McGuinness, 2003)

■ $\chi(G^2) \leq \frac{9}{5} \Delta(G) + 1$ (for $\Delta(G) \geq 47$)
(Borodin, Broersma, Glebov & vdH, 2001)

■ $\chi(G^2) \leq \frac{5}{3} \Delta(G) + 24$ (for $\Delta(G) \geq 241$)
(Molloy & Salavatipour, 2005)

Towards Wegner's Conjecture

Theorem (Havet, vdH, McDiarmid & Reed, 2008+)

■ G planar $\implies \chi(G^2) \leq (3/2 + \epsilon) \Delta(G)$
($\epsilon \downarrow 0$ for $\Delta(G) \rightarrow \infty$)

Theorem (Amini, Esperet & vdH, 2013)

■ G embeddable on a fixed surface S

■ $\implies \chi(G^2) \leq (3/2 + \epsilon) \Delta(G)$ ($\Delta(G) \rightarrow \infty$)

■ \implies clique number $\omega(G^2) \leq 3/2 \Delta(G) + C$

What about distances larger than 2 ?

easy upper bound

$$\blacksquare \chi(G^d) \leq 1 + \sum_{i=0}^{d-1} \Delta(G) (\Delta(G) - 1)^i = \Omega(\Delta(G)^d)$$

Theorem (Agnarsson & Halldórsson, 2003)

$$\blacksquare G \text{ } k\text{-degenerate} \implies \chi(G^d) \leq c_{k,d} \Delta(G)^{\lfloor d/2 \rfloor}$$

Colouring the cube of planar graphs

- so there is some constant c_3 such that:

$$G \text{ planar} \implies \chi(G^3) \leq c_3 \Delta(G) + C$$

- but what is the best c_3 ?

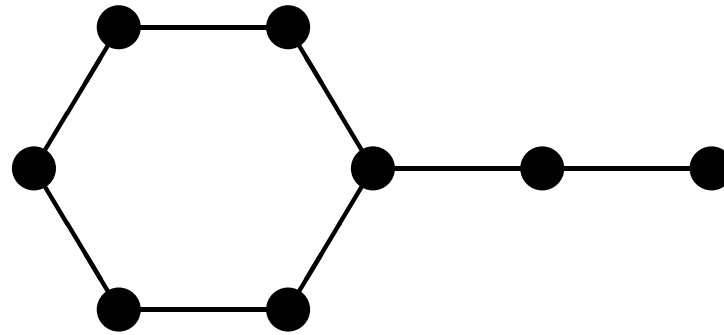
- we only know: $9/2 \leq c_3 \leq 45$

- and what about distances $d > 3$?

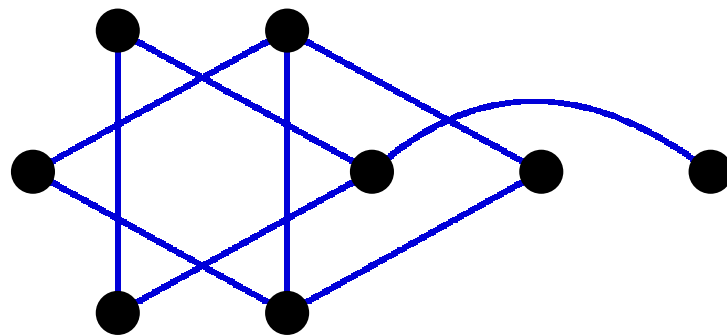
A variant with exact distances

- suppose we only want vertices **at distance exactly d** to have **different colours**
- can be modelled using the **exact distance graph $G^{[d]}$** :
 - same vertex set as G
 - edges between vertices with **distance exactly d** in G

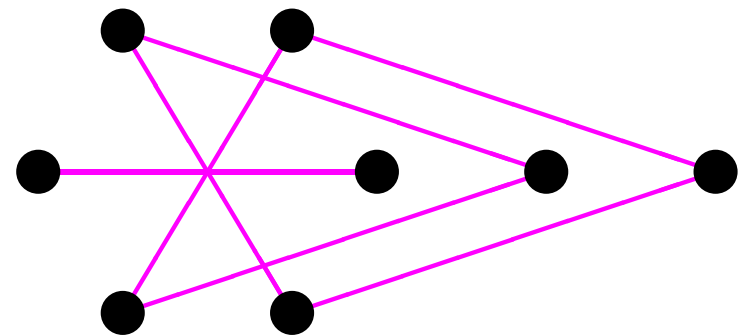
Exact distance graphs



G



$G^{[2]}$



$G^{[3]}$

Colouring at an exact distance

- for even d , for most classes of graphs the chromatic number $\chi(G^{[\#d]})$ is not bounded
- but for odd d , the situation is quite different

Theorem (Nešetřil & Ossona de Mendez, 2008)

- \mathcal{G} a class of graphs with bounded expansion, d odd
 \implies there exists a constant $N_{\mathcal{G},d}$ such that:
for all $G \in \mathcal{G}$: $\chi(G^{[\#d]}) \leq N_{\mathcal{G},d}$

A very, very special case

- the class of **planar graphs** has **bounded expansion**

SO ...

A very, very special case

Corollary

- there exists a constant C such that

$$G \text{ planar} \implies \chi(G^{[\#3]}) \leq C$$

- ■ proof of Nešetřil & Ossona de Mendez is long, complicated, and gives little idea what is going on
- $G^{[\#3]}$ can be very dense for planar G
- there is no bound on the list-chromatic number of $G^{[\#3]}$
- until recently, best known bounds on C :

$$6 \leq C \leq 5 \cdot 2^{10,241}$$

A very, very simple result

Theorem (vdH, Kierstead & Quiroz, 2016)

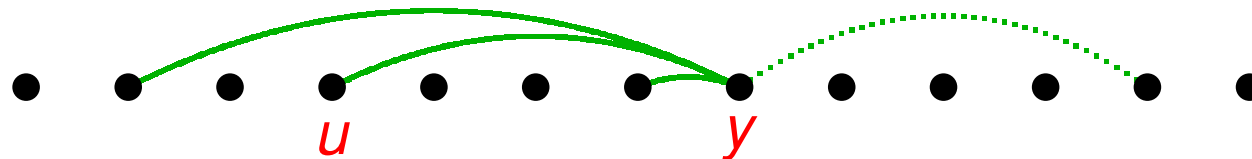
- d odd, then for every graph G :

$$\chi(G^{[\#d]}) \leq \text{wcol}_{2d-1}(G)$$

- the **weak d -colouring number** $\text{wcol}_d(G)$ is a generalisation of **degeneracy**

The normal colouring number

- let L be a linear ordering of the vertices of a graph G
- for a vertex $y \in V(G)$,
let $S(L, y)$ be the set of neighbours u of y with $u <_L y$



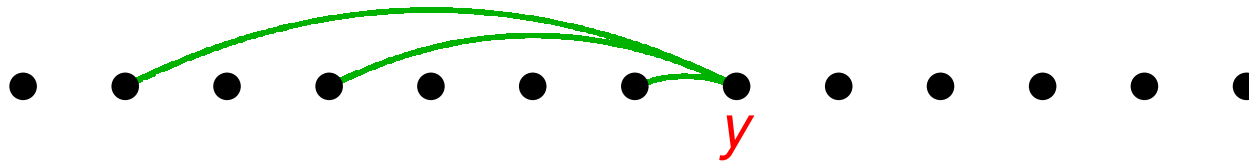
- then the colouring number $\text{col}(G)$ is defined as

$$\text{col}(G) = \min_L \max_{y \in V(G)} |S(L, y)| + 1$$

- note: G k -degenerate $\implies \text{col}(G) \leq k + 1$

Generalising the colouring number

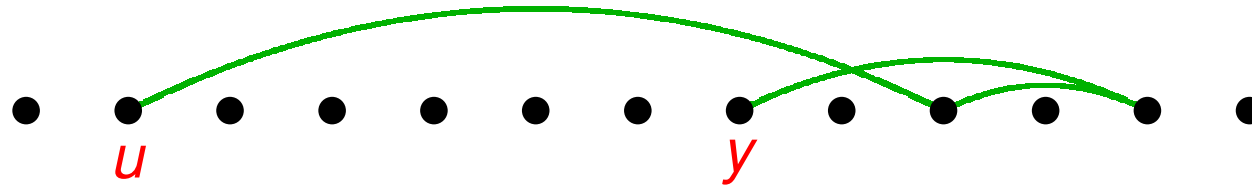
- the set $S(L, y)$ can be defined as
“vertices $u <_L y$ for which there is a uy -path of length l ”



- what would happen if we allow longer paths ?

The strong colouring number

- a **strong uy -path** has all internal vertices larger than y

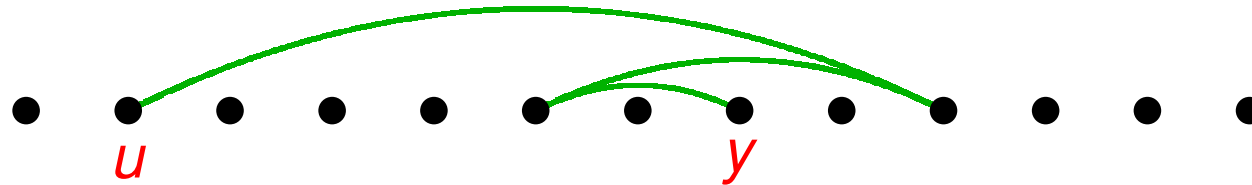


- let $S_d(L, y)$ be the set of vertices $u <_L y$ for which there exists a strong uy -path with length at most d
- then the **strong d -colouring number $scol_d(G)$** is defined as

$$scol_d(G) = \min_L \max_{y \in V(G)} |S_d(L, y)| + 1$$

The weak colouring number

- a **weak uy -path** has all internal vertices larger than u



- let $W_d(L, y)$ be the set of vertices $u <_L y$ for which there exists a weak uy -path of length at most d
- then the **weak d -colouring number** $\text{scol}_d(G)$ is defined as

$$\text{wcol}_d(G) = \min_L \max_{y \in V(G)} |W_d(L, y)| + 1$$

Some basic facts of these generalised colouring numbers

- by definition: $\text{col}(G) = \text{scol}_1(G) = \text{wcol}_1(G)$
- ■ obviously: $\text{scol}_d(G) \leq \text{wcol}_d(G)$
 - in fact, also: $\text{wcol}_d(G) \leq (\text{scol}_d(G))^d$
- hence:
if **one** of scol_d , wcol_d , is **bounded** on some class of graphs,
then **the other one** is also **bounded** on that class

Not so basic fact

- again straightforward from the definition:

$$\text{scol}_1(G) \leq \text{scol}_2(G) \leq \text{scol}_3(G) \leq \dots \leq \text{scol}_\infty(G)$$

(where the “ ∞ ” means “any length strong path allowed”)

- Property (Grohe et al., 2014)

$$\text{scol}_\infty(G) = \text{treewidth}(G) + 1$$

Back to classes with bounded expansion

Definition (Nešetřil & Ossona de Mendez)

- a class of graphs \mathcal{G} has **bounded expansion** if there exist constants c_1, c_2, \dots , such that for all $G \in \mathcal{G}$ and for all $d = 1, 2, \dots$ we have:
 - for all minors H of G formed by contracting connected subgraphs with radius at most d :

$$|E(H)| \leq c_d \cdot |V(H)|$$

- generalises classes with “bounded treewidth”, “bounded genus”, “forbidden minors”, “bounded cop number”, etc., etc.

Back to classes with bounded expansion

Equivalent definitions (Zhu, 2009)

a class of graphs \mathcal{G} has **bounded expansion** if

- there exist constants c'_1, c'_2, \dots such that

- for all $G \in \mathcal{G}$ and for all $d = 1, 2, \dots$:

$$\text{scol}_d(G) \leq c'_d$$

or

- there exist constants c''_1, c''_2, \dots such that

- for all $G \in \mathcal{G}$ and for all $d = 1, 2, \dots$:

$$\text{wcol}_d(G) \leq c''_d$$

Back to colouring exact distance graphs

time to prove that for any graph G :

$$\chi(G^{[3]}) \leq \text{wcol}_5(G)$$

- recall the definition of $\text{wcol}_5(G)$:

$$\text{wcol}_5(G) = \min_L \max_{y \in V(G)} |W_5(L, y)| + 1$$

- so let's choose an ordering L_5 of the vertices such that

$$\text{for all } y: |W_5(L_5, y)| + 1 \leq \text{wcol}_5(G)$$

Back to colouring exact distance graphs

- so let's choose an **ordering** L_5 of the vertices such that
for all y : $|W_5(L_5, y)| + 1 \leq wcol_5(G)$
- stage 1
going along the **ordering** L_5 , give every vertex y
a colour $c(y)$ different from the vertices in $W_5(L_5, y)$
 - requires **at most** $wcol_5(G)$ colours

Back to colouring exact distance graphs

- stage 1

going along the ordering L_5 , give every vertex y a colour $c(y)$ different from the vertices in $W_5(L_5, y)$

- stage 2

- for a vertex y , define $N[y] = N(y) \cup \{y\}$

and let $\mu(y)$ be the left-most vertex in $N[y]$

(according to the ordering L_5)

- give every vertex y the colour $C(y) = c(\mu(y))$

- Claim: the colouring C is a proper colouring of $G^{[\#3]}$ \square

What can we say about $wcol_d(G)$?

Theorem (vdH, Kierstead & Quiroz, 2016)

- d odd, then for every graph G :

$$\chi(G^{[\#d]}) \leq wcol_{2d-1}(G)$$

Theorem (vdH, Ossona de Mendez, Quiroz, Rabinovich & Siebertz, 2016)

- bounds on $scol_d(G)$ and $wcol_d(G)$ for all kinds of graphs (bounded treewidth, bounded genus, forbidden minor, etc.)

- in particular:

- G planar $\implies wcol_d(G) \leq \binom{d+2}{2} \cdot (2d+1)$

Bounds on $\chi(G^{[\#3]})$ for planar graphs

Corollary (vdH, Kierstead & Quiroz, 2016)

- G planar $\implies \chi(G^{[\#3]}) \leq \text{wcol}_5(G) \leq 231$
- by being a bit more careful, we can prove:
 G planar $\implies \chi(G^{[\#3]}) \leq 143$
- we also constructed a planar H with $\chi(H^{[\#3]}) = 7$

Some final bits and bobs

since a couple of weeks we know

Theorem

(Bousquet, Esperet, Harutyunyan, de Joannis de Verclos, Pastor)

$$\blacksquare \max \{ \chi(G^{[\#d]}) \mid G \text{ planar} \} \longrightarrow \infty \text{ if } d \longrightarrow \infty$$

Some final bits and bobs

Theorem (vdH, Kierstead & Quiroz, 2016)

- d odd, then for every graph G :

$$\chi(G^{[\#d]}) \leq \text{wcol}_{2d-1}(G)$$

- d even, then for every graph G :

$$\chi(G^{[\#d]}) \leq \text{wcol}_{2d}(G) \cdot \Delta(G)$$

Some final bits and bobs

Theorem (vdH, Kierstead & Quiroz, 2016)

- d even, then for every graph G :

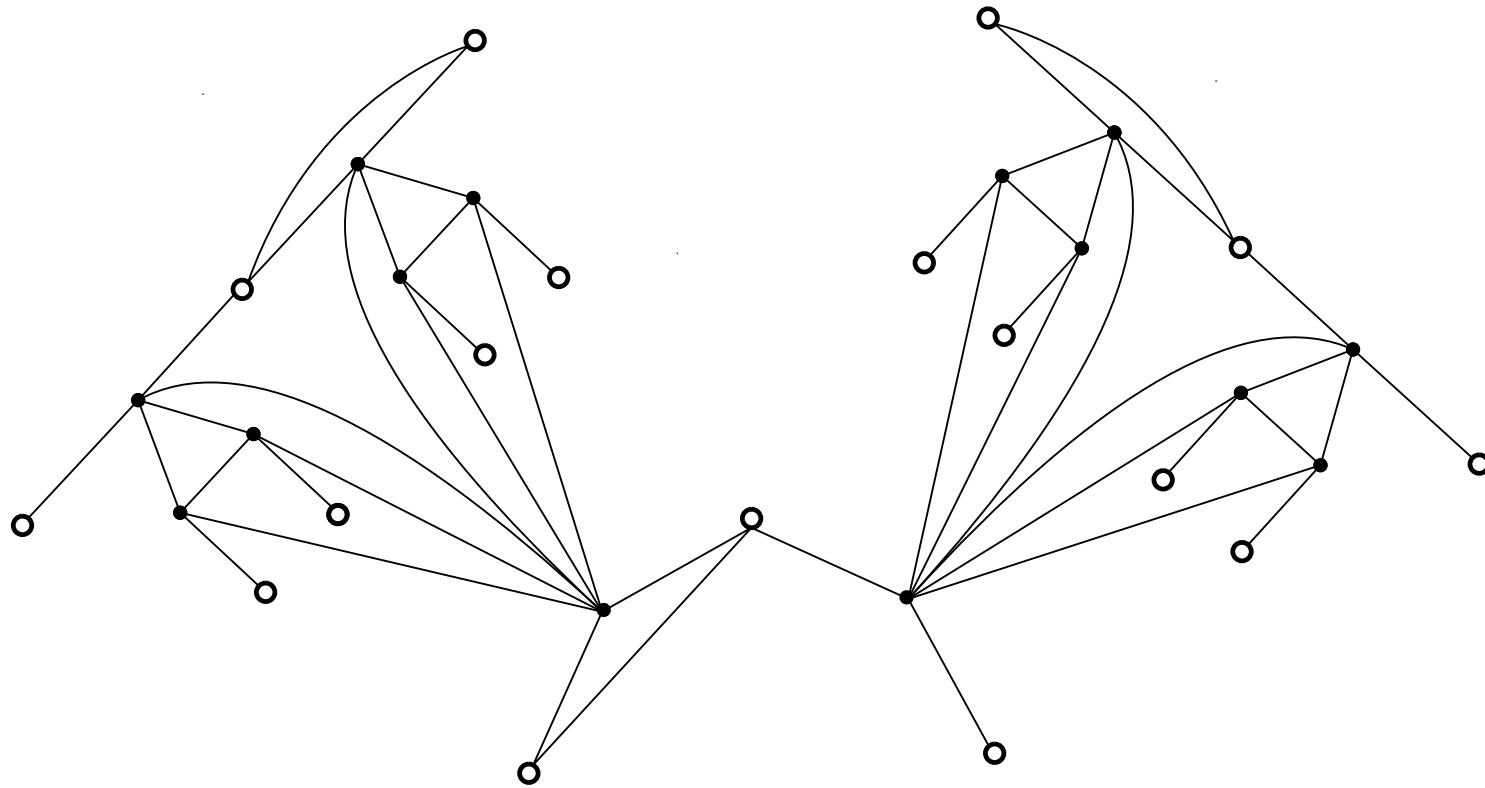
$$\chi(G^{[\#d]}) \leq \text{wcol}_{2d}(G) \cdot \Delta(G)$$

Corollary

- \mathcal{G} a class of graphs with bounded expansion,
 d even

\implies there exists a constant $N'_{\mathcal{G},d}$ such that:

$$\text{for all } G \in \mathcal{G} : \chi(G^{[\#d]}) \leq N'_{\mathcal{G},d} \cdot \Delta(G)$$



Thanks for your attention!