Improper Colourings inspired by Hadwiger's Conjecture

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Graph minors

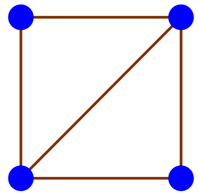
- a graph H is a **minor** of a graph G, if H can be obtained from G by a series of:
  - vertex deletions;
  - edge deletions;
  - edge contractions:

Graph minors – more intuitive

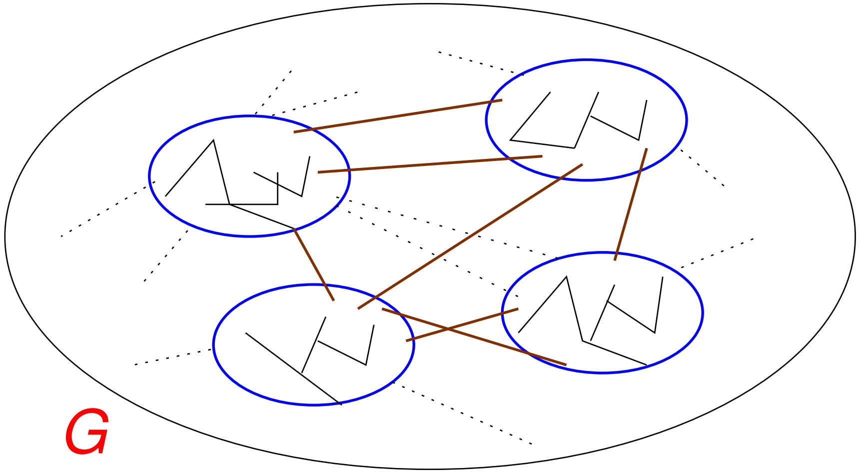
- a graph H is a **minor** of a graph G if:

  for $V(H) = \{v_1, \ldots, v_k\}$, there exist connected, disjoint subgraphs $H_1, \ldots, H_k$ of G such that:

  - if $v_i v_j \in E(H)$, then there is at least one edge in G between H_i and H_j



H



G

Graph colouring

- a **colouring** of a graph means colouring the vertices
- **proper colouring**: adjacent vertices have different colours

- recurring question in graph theory:

  what structural properties of a graph
  - allow proper colourings with few colours ?
  - force all proper colourings to use many colours ?

Hadwiger's Conjecture

Conjecture (Hadwiger, 1943)

- a graph G needs at least t colours for a proper colouring

  $\Longrightarrow$ G has the complete graph K_t as a minor

the contrapositive is probably more intuitive:

- G has no K_t -minor

  $\stackrel{?}{\Longrightarrow}$ G has a proper $(t-1)$ -colouring

Hadwiger's Conjecture – what is known

- no K_t -minor $\stackrel{?}{\Longrightarrow}$ G has a proper $(t-1)$ -colouring

known

- nothing to prove for $t = 1, 2$
- easy for $t = 3$
- not too hard for $t = 4$ (Hadwiger, 1943; Dirac, 1952)
- case $t = 5$ is equivalent to the Four Colour Theorem (Wagner, 1937)
- true for $t = 6$ (Robertson, Seymour & Thomas, 1993)

How many colours do we need ?

Theorem (Kostochka, 1984; Thomason, 1984)

- G has no K_t-minor

  $\Longrightarrow$ G has a vertex with degree at most $c\,t\sqrt{\log t}$

Corollary

- G has no K_t-minor

  $\Longrightarrow$ G has a proper colouring with $c\,t\sqrt{\log t}$ colours

Improper colourings

- what if we weaken the requirement on the colouring?
  - in a proper colouring:

    the collection of vertices with the same colour
    is just a collection of isolated vertices

  - we could be happy with:

    the collection of vertices with the same colour
    is just a subgraph with a "simple" structure

- **monochromatic subgraph**:
  subgraph formed by vertices with the same colour

Improper colourings – small monochromatic degree

Theorem (Edwards, Kang, Kim, Oum & Seymour, 2015)

- G has no K_t-minor $\Longrightarrow$

  G can be coloured with $t-1$ colours such that each monochromatic subgraph has degree at most $c'\,t^2 \log t$

- the bound $t-1$ on the number of colours is best possible:
  - there exists a class of graphs without K_t-minor,

    but where you can't bound the degree of monochromatic subgraphs when using $t-2$ colours only

Improper colourings – small monochromatic degree

Theorem (Edwards, Kang, Kim, Oum & Seymour, 2015)

- G has no K_t -minor $\Longrightarrow$

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Theorem (vdH & Wood, 2017)

- G has no K_t -minor $\Longrightarrow$

  G can be coloured with $t-1$ colours such that each monochromatic subgraph has degree at most $t-2$

no K_t -minor $\overset{?}{\Longrightarrow}$ proper $(t-1)$ -colouring

Improper colourings – small monochromatic components

Theorem (Kawarabayashi & Mohar, 2007)

- G has no K_t-minor $\Longrightarrow$

  G can be coloured with $\lceil 15\frac{1}{2}\,t \rceil$ colours such that each monochromatic components has at most $f_1(t)$ vertices

- improved to
  - $\lceil 3\frac{1}{2}\,t - 1\frac{1}{2} \rceil$ colours; $f_2(t)$ vertices (Wood, 2010 (?))
  - $3(t-1)$ colours; $f_3(t)$ vertices (Liu & Oum, 2015)
  - $2(t-1)$ colours; $f_4(t)$ vertices (Norin, 2015; unpubl.)

- all use Robertson & Seymour Graph Minor Structure Thm., or worse ...

Improper colourings – small monochromatic components

Theorem (vdH & Wood, 2017)

- G has no K_t -minor $\Longrightarrow$

  G can be coloured with $2(t-1)$ colours such that each monochromatic component has at most $\lceil \frac{1}{2}(t-2) \rceil$ vertices

note

- G has no K_t -minor $\Longrightarrow$

  at least $t-1$ colours are needed to guarantee monochromatic components of bounded size

  (same examples as for small monochromatic degree)

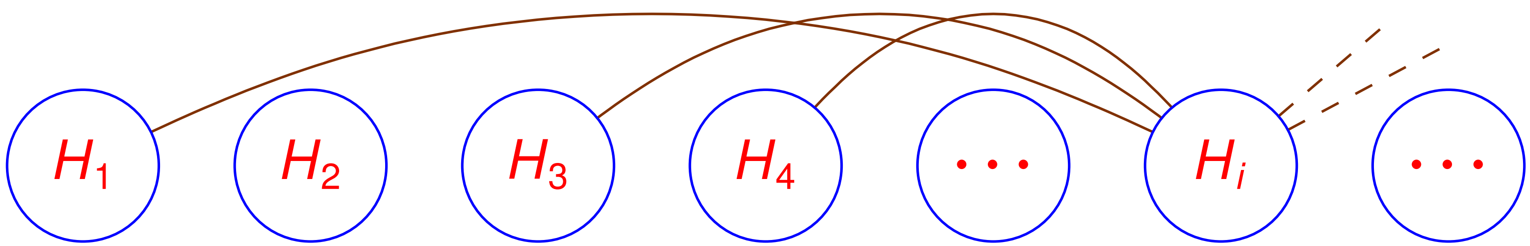
A simple decomposition theorem for K_t-minor-free graphs

Theorem (vdH & Wood, 2017)

- G has no K_t -minor $\Longrightarrow$

  G has a partition into subgraphs $H_1, \ldots, H_\ell$ such that

  - global structure: each H_i is adjacent to at most $t-2$ of the earlier subgraphs $H_1, \ldots, H_{i-1}$



A simple decomposition theorem for K_t-minor-free graphs

Theorem (vdH & Wood, 2017)

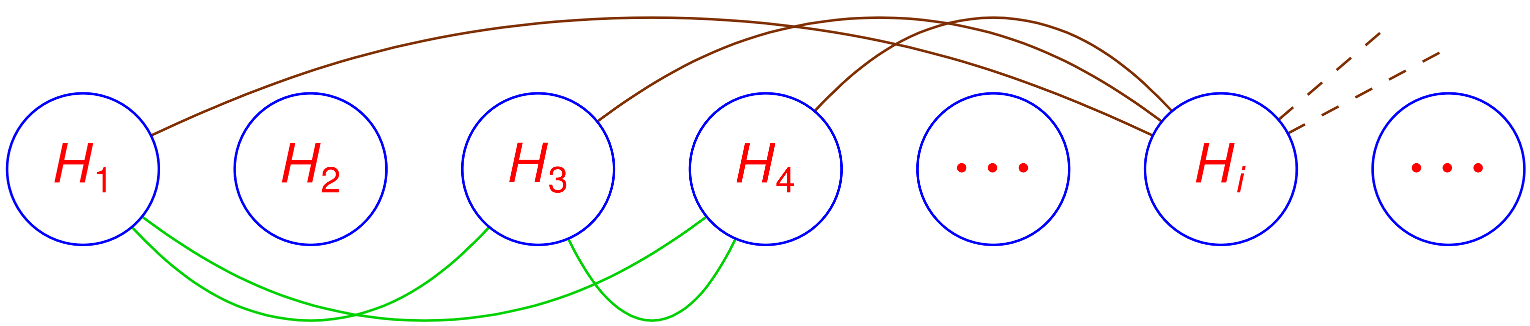
- G has no K_t-minor $\Longrightarrow$

  G has a partition into subgraphs $H_1, \ldots, H_\ell$ such that

  - global structure: each H_i is adjacent to at most $t-2$ of the earlier subgraphs $H_1, \ldots, H_{i-1}$
  - local structure:
    - each H_i has maximum degree at most $t-2$
    - each H_i can be coloured with 2 colours such that each monochromatic component of H_i has at most $\left\lceil \frac{1}{2}(t-2) \right\rceil$ vertices

The global structure we actually prove

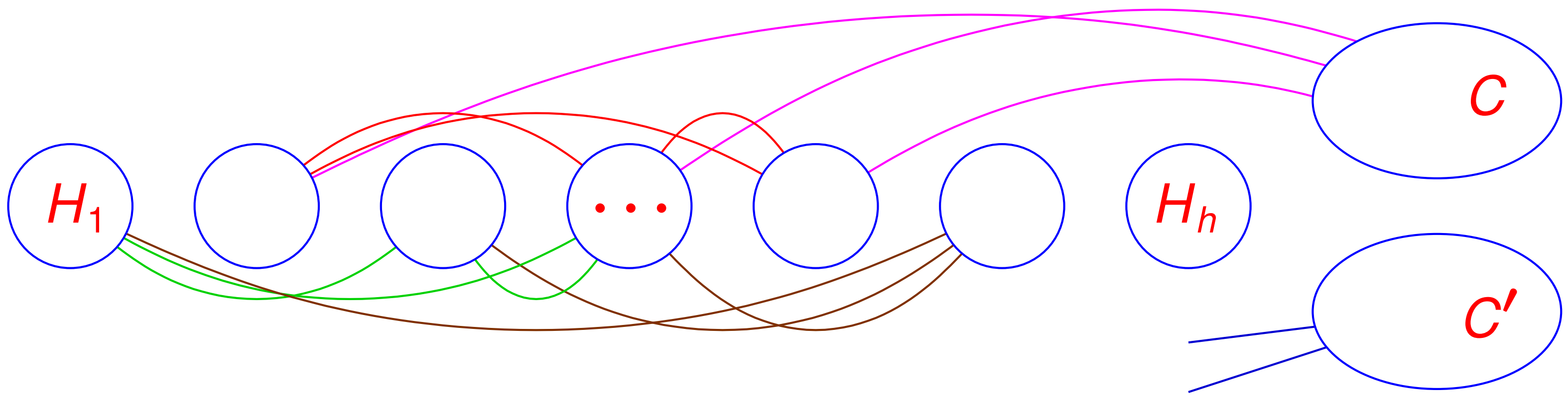
- G any graph $\Longrightarrow$
  we can construct (in many ways) a partition of G into induced subgraphs $H_1, \ldots, H_\ell$ such that:
  - each H_i is connected
  - each H_i is adjacent to k subgraphs $H_{i_1}, \ldots, H_{i_k}$ the earlier subgraphs $H_1, \ldots, H_{i-1}$
  - for each H_i, the adjacent subgraphs $H_{i_1}, \ldots, H_{i_k}$ are pairwise adjacent as well



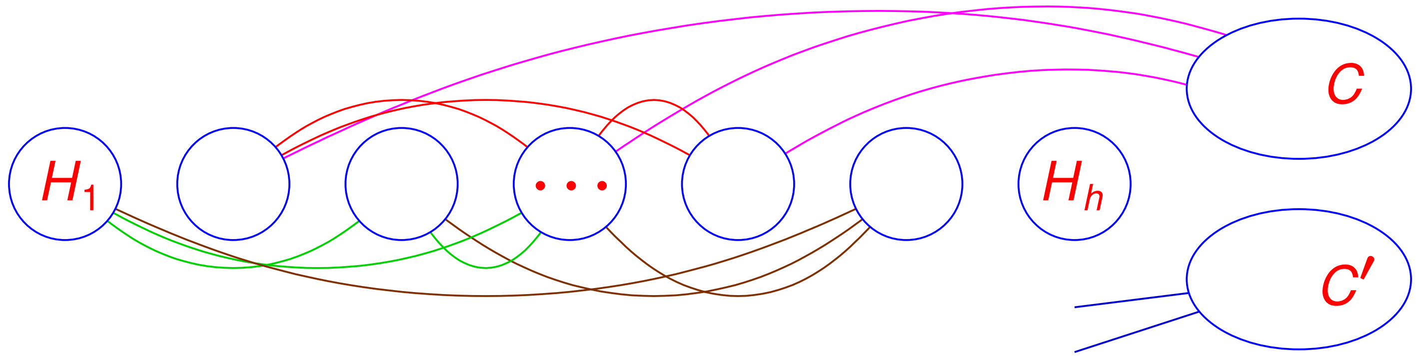
The global structure – proof

- we will construct the H_i one by one such that once $H_1, \ldots, H_h$ is constructed:
  - each H_i, $i \leq h$, satisfies the requirements
  - each component C of $G - \big(V(H_1) \cup \cdots \cup V(H_h)\big)$ satisfies:

    if C is adjacent to $H_{i_1}, \ldots, H_{i_k}$ from $H_1, \ldots, H_{i-1}$, then $H_{i_1}, \ldots, H_{i_k}$ are pairwise adjacent as well

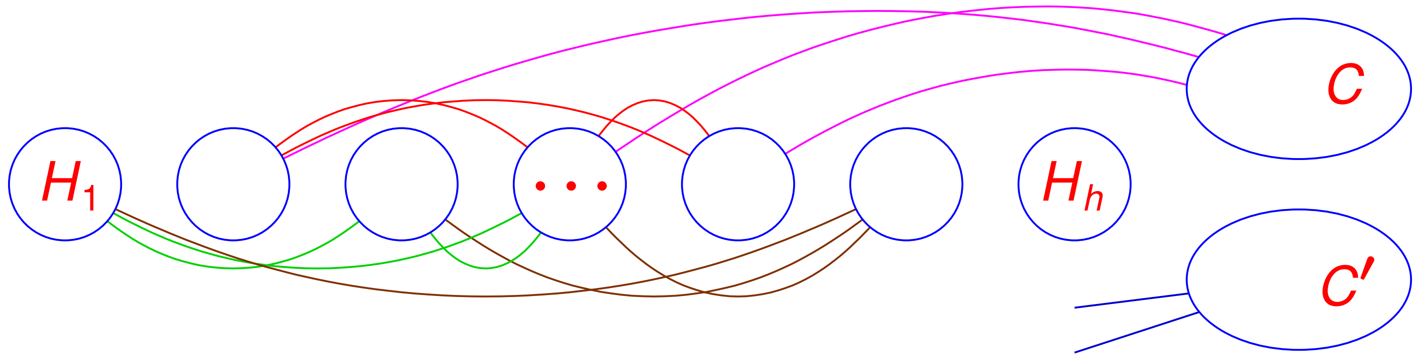


The global structure – proof



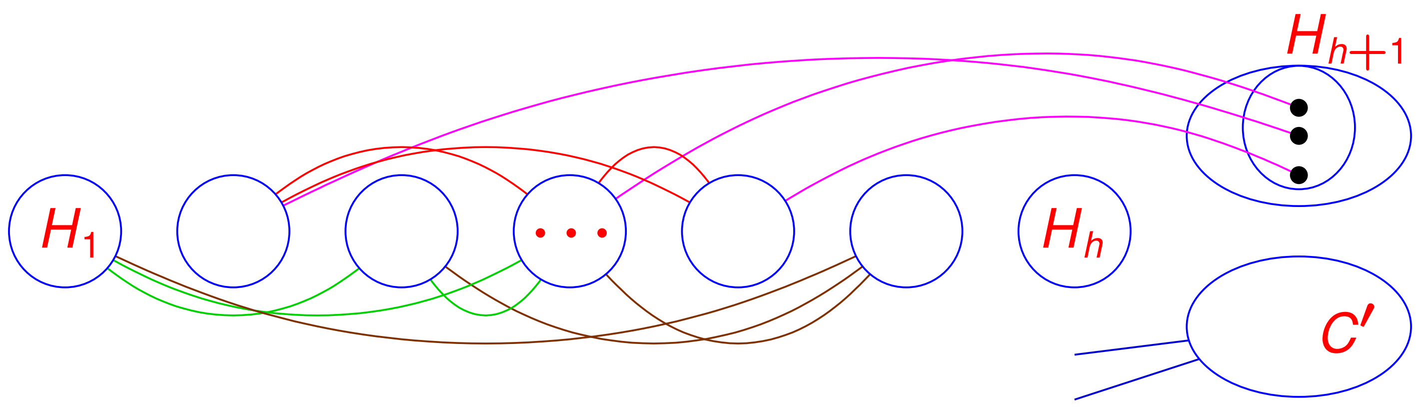
- start with H_1 any connected, induced subgraph of G
  (good choice for later: $V(H_1) = \{v\}$ for some $v \in V(G)$)
  - all requirements are trivially satisfied

The global structure – proof



- suppose $H_1, \ldots, H_h$ are already constructed and C is some component of $G - \big(V(H_1) \cup \cdots \cup V(H_h)\big)$
  - so C is adjacent to $H_{i_1}, \ldots, H_{i_k}$, which are also pairwise adjacent
- for each H_{i_ℓ}, choose $a_{i_\ell} \in V(C)$ adjacent to H_{i_ℓ}
- now choose H_{h+1} a connected, induced subgraph of C containing all $a_{i_1}, \ldots, a_{i_k}$

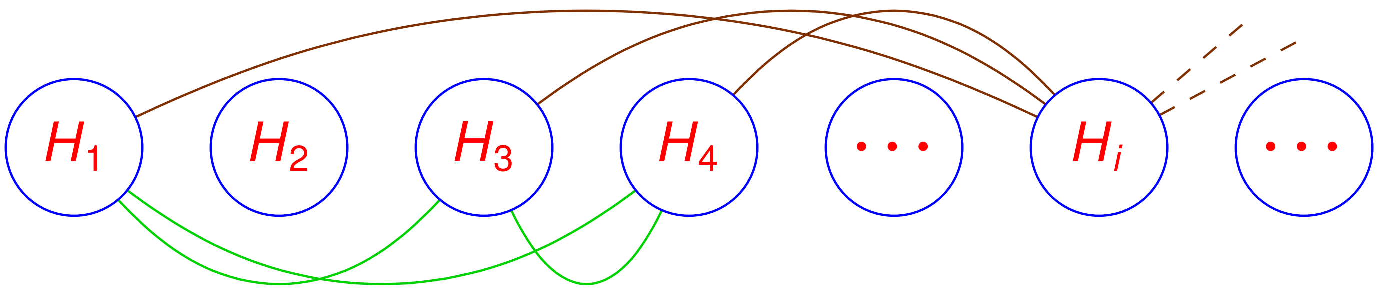
The global structure – proof



- H_{h+1} is still adjacent to $H_{i_1}, \ldots, H_{i_k}$
- a component of $G - \big(V(H_1) \cup \cdots \cup V(H_h) \cup V(H_{h+1})\big)$
  - is either a component of $G - \big(V(H_1) \cup \cdots \cup V(H_h)\big)$ and hence still satisfies the requirements
  - or it is a component of $C - V(H_{h+1})$, hence it is adjacent to H_{h+1} and some of $H_{i_1}, \ldots, H_{i_k}$, which are all pairwise adjacent

The global structure for K_t-minor-free graphs

- we can construct (in many ways) a partition of any graph G into induced, connected subgraphs $H_1, \ldots, H_\ell$ such that:
  - each H_i is adjacent to k subgraphs $H_{i_1}, \ldots, H_{i_k}$ from $H_1, \ldots, H_{i-1}$, which are pairwise adjacent as well



- G has no K_t-minor $\Longrightarrow$

  for each H_i we must have $k \le t-2$

The local structure – inside the H_i

- so how can we choose the H_i so that they satisfy:
  - small degree and
  - 2-colourable with small mono. components
- each H_i was chosen as some induced subgraph of some connected subgraph C, such that:
  - H_i is connected
  - H_i contains some set $A = \{a_{i_1}, \ldots, a_{i_k}\}$
- idea:
  choose H_i the smallest subgraph with those properties

The local structure – inside the H_i

Lemma

- C a connected graph, $A \subseteq V(C)$
  H a minimal, induced, connected subgraph of C,
  such that H contains all of A

- then H satisfies:
  - every vertex in H has degree at most $|A|$ in H
  - every vertex not in A is a cut-vertex of H
    - easy corollary: there is a 2-colouring of H with monochromatic components of size at most $\lceil \frac{1}{2}|A| \rceil$

Our decomposition theorem again

Theorem (vdH & Wood, 2017)

- G has no K_t-minor $\Longrightarrow$

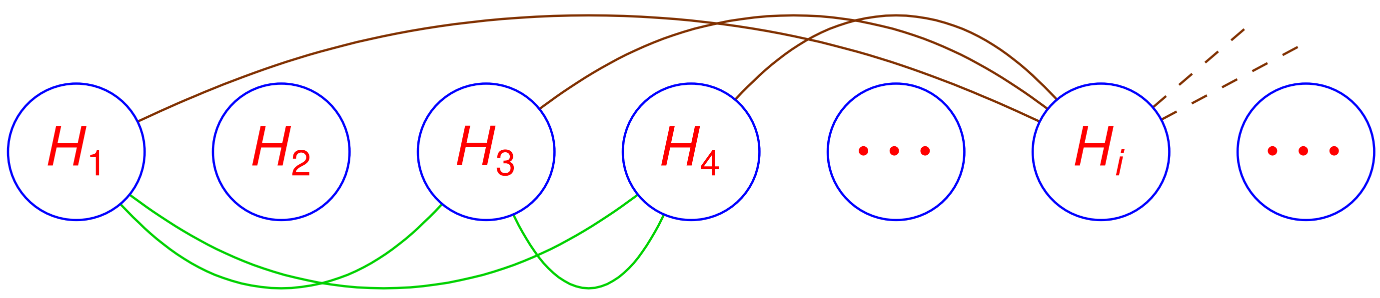
  G has a partition into subgraphs $H_1, \ldots, H_\ell$ such that

  - global structure: each H_i is adjacent to at most $t-2$ of the earlier subgraphs $H_1, \ldots, H_{i-1}$
  - local structure:
    - each H_i has maximum degree at most $t-2$
    - each H_i can be coloured with 2 colours such that each monochromatic component of H_i has at most $\left\lceil \frac{1}{2}(t-2) \right\rceil$ vertices
  - some more properties

With only a little bit extra work ...

Theorem

- G has no K_t-minor $\Longrightarrow$ G has a partition into connected subgraphs $H_1, \ldots, H_\ell$ such that
  - contracting all H_i to single vertices gives a chordal graph with treewidth at most $t-2$
  - each H_i has treewidth at most $t-3$

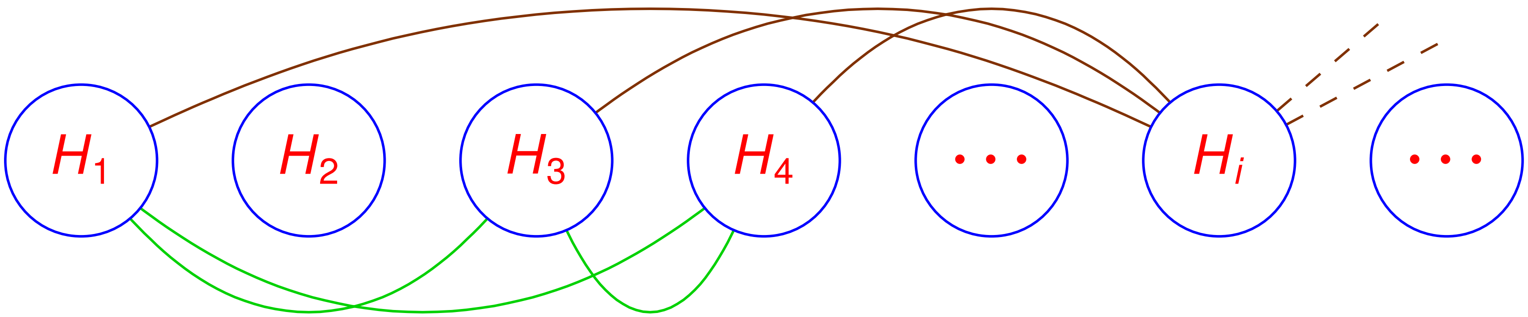


A similar result for $K_{3,s}$-minor-free graphs

- we can prove a similar decomposition theorem for $K_{3,s}$-minor-free graphs

Corollary

- G has no $K_{3,s}$-minor $\Longrightarrow$
  - G can be coloured with 3 colours such that each monochromatic subgraph has degree at most $4s$
  - G can be coloured with 6 colours such that each monochromatic component has at most $2s$ vertices



That's all folks! – Thanks for listening.