# Strategic Characterization of the Index of an Equilibrium 

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## Battle of the sexes



## Battle of the sexes with in-laws



## Battle of the sexes with in-laws



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## Make NE unique by adding strategies

Given:
nondegenerate (A,B), Nash equilibrium (NE) ( $\mathrm{x}, \mathrm{y}$ ).
Question:
Is there a game G extending ( $\mathrm{A}, \mathrm{B}$ ) by adding strategies so that ( $x, y$ ) is the unique NE of $G$ ?
e.g.: G obtained from ( $\mathrm{A}, \mathrm{B}$ ) by adding columns, $(x, y)$ becoming ( $x,[y, 0,0, \ldots, 0]$ ) for $G$.

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- can be done for pure-strategy NE
- but not for mixed NE of "battle of the sexes"


## Strategic characterization of the index

We will show a conjecture by Josef Hofbauer:

## Theorem:

For nondegenerate (A,B), Nash equilibrium ( $\mathrm{x}, \mathrm{y}$ ):

$$
\text { index }(x, y)=+1
$$

$\Leftrightarrow \quad \exists$ game G extending (A,B)
so that ( $x, y$ ) is the unique equilibrium of $G$.
for $m \times n$ game:
G obtained from ( $\mathrm{A}, \mathrm{B}$ ) by adding 3 m columns.

## Sub-matrices of equilibrium supports

Given: nondegenerate $(A, B), A>0, B>0$, Nash equilibrium ( $\mathrm{x}, \mathrm{y}$ ).

$$
\begin{aligned}
& A=\left(a_{i j}\right), B=\left(b_{i j}\right) \\
& A_{x y}=\left(a_{i j}\right) i \in \operatorname{supp}(x), j \in \operatorname{supp}(y) \\
& B_{x y}=\left(b_{i j}\right) \quad i \in \operatorname{supp}(x), j \in \operatorname{supp}(y)
\end{aligned}
$$

Axy, $B_{x y}$ have full rank |supp(x)|, nonzero determinants.

## Index of an equilibrium (Shapley 1974)

Given: nondegenerate game $(A, B), \quad A>0, B>0$,
Nash equilibrium ( $x, y$ ).
index $(x, y)=(-1)|\operatorname{supp}(x)|+1$ sign $\operatorname{det}\left(A_{x y} B_{x y}\right)$

$$
\in\{+1,-1\}
$$

## Properties of the index

- independent of
- positive constant added to all payoffs
- order of pure strategies
- pure strategy payoffs outside equilibrium support
- pure-strategy equilibria have index +1
- sum of indices over all equilibria is +1
- the two endpoints of any Lemke-Howson path are equilibria of opposite index.


## Definition of symmetric index

Given: nondegenerate symmetric game $\left(B^{\top}, B\right), \quad B>0$, SNE (symmetric Nash equilibrium) ( $x, x$ ).
symmetric index $(x, x)=(-1)^{|\operatorname{supp}(x)|+1} \operatorname{sign} \operatorname{det}\left(B_{x x}\right)$

$$
\in\{+1,-1\}
$$

## Symmetric NE of symmetric games



## best-reply regions



## symmetric equilibria



## symmetric equilibria



## symmetric equilibria



## symmetric equilibria



## Only ■dynamically stable, $\triangleleft$ not



## best-reply regions



## best-reply regions



## best-reply regions



## symmetric Lemke-Howson paths



## symmetric Lemke-Howson paths



## Points instead of best-reply regions

$\left.B=$| 3 | 4 | 5 |
| :--- | :--- | :--- |
| 1 | 3 | 4 |
| 3 | 2 | 1 | \right\rvert\,

## Mixed strategies of player 1



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## Non-negative convex hull



## Use negative unit vectors instead



## Use negative unit vectors instead



## Project to get Schlegel diagram



## Project to get Schlegel diagram



Project to get Schlegel diagram

$$
\left.B=\begin{array}{|lll}
3 & 4 & 5 \\
1 & 3 & 4 \\
3 & 2 & 1
\end{array} \right\rvert\,
$$



## Example: m=3 strategies for player 1

$$
\left.\begin{array}{c}
\left.A=\begin{array}{|cccc}
0 & 10 & 0 & 10 \\
10 & 0 & 0 & 0 \\
8 & 0 & 10 & 8
\end{array}\right] \\
\left.B=\begin{array}{|cccc}
0 & 10 & 0 & -10 \\
0 & 0 & 10 & 8 \\
10 & 0 & 0 & 8
\end{array}\right] \\
P^{\Delta}=\operatorname{conv}\left(\begin{array}{ccccc}
\mathbf{1} & \mathbf{2} & \mathbf{3} & \mathbf{4} & \mathbf{5} \\
-\mathrm{M} & \mathbf{6} & \mathbf{7} \\
0 & -\mathrm{M} & 0 & 0 & 10 \\
0 & 0 & -10 \\
0 & 0 & -\mathrm{M} & 10 & 0 \\
10 & 8 \\
\hline
\end{array}\right.
\end{array}\right), \quad \mathrm{M} \text { large }, \quad . \quad .
$$




## Schlegel diagram for $\mathbf{P}^{\Delta}$



## subdivide facet: player 1's best responses



## Facets of $\mathbf{P}^{\Delta}=$ potential equilibrium strategies of player 1

## Each facet of $\mathrm{P}^{\Delta}$

$=$ simplex spanned by m columns of $[-\mathrm{M} \mathrm{Im}, \mathrm{B}]$
Normal vector of facet
$=$ mixed strategy of player 1
m columns
$=$ best responses of player 2 or unused strategies of player 1

# Subdivide facets into special best response regions of player 1 

$\left.\begin{array}{l|ccccccc}\text { use } & \mathbf{1} & \mathbf{2} & \mathbf{3} & \mathbf{4} & \mathbf{5} & \mathbf{6} & \mathbf{7} \\ \text { matrix } \mathrm{B}: & -\mathrm{M} & 0 & 0 & 0 & 10 & 0 & -10 \\ \mathrm{P}^{\Delta}=\operatorname{conv}( & -M & 0 & 0 & 0 & 10 & 8 \\ 0 & 0 & -\mathrm{M} & 10 & 0 & 0 & 8\end{array}\right)$

|  | $\mathbf{1}$ | $\mathbf{2}$ | $\mathbf{3}$ | $\mathbf{4}$ | $\mathbf{5}$ | $\mathbf{6}$ | $\mathbf{7}$ |
| :--- | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| use | $\mathbf{1}$ | 0 | 0 | 0 | 10 | 0 | 10 |
| 0 | 1 | 0 | 10 | 0 | 0 | 0 |  |
| 0 | 0 | 1 | 8 | 0 | 10 | 8 |  |

## Subdivide facets into special best response regions of player 1

$$
\left.\mathrm{P}^{\Delta}=\operatorname{conv}\left(\begin{array}{ccc|ccc|c}
\mathbf{1} & \mathbf{2} & \mathbf{3} & \mathbf{4} & \mathbf{5} & \mathbf{6} & \mathbf{7} \\
-\mathrm{M} & 0 & 0 & 0 & 10 & 0 & -10 \\
0 & -\mathrm{M} & 0 & 0 & 0 & 10 & 8 \\
0 & 0 & -\mathrm{M} & 10 & 0 & 0 & 8
\end{array}\right]\right) \quad \text { facet }
$$


subdivide into
best response
regions

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$$
\mathrm{P}^{\Delta}=\operatorname{conv}\left(\begin{array}{c|cc|ccc|c|}
\mathbf{1} & \mathbf{2} & \mathbf{3} & \mathbf{4} & \mathbf{5} & \mathbf{6} & \mathbf{7} \\
-\mathrm{M} & 0 & 0 & 0 & 10 & 0 & -10 \\
\mathbf{0} & -\mathrm{M} & 0 & 0 & 0 & 10 & 8 \\
\mathbf{0} & 0 & -\mathrm{M} & 10 & 0 & 0 & 8 \\
\hline
\end{array}\right) \quad \begin{aligned}
& \text { facet with } \\
& \text { unplayed } \\
& \text { strategy: }
\end{aligned}
$$

| $\mathbf{1}$ | $\mathbf{2}$ | $\mathbf{3}$ | $\mathbf{4}$ | $\mathbf{5}$ | $\mathbf{6}$ | $\mathbf{7}$ |
| :--- | :--- | :--- | :---: | :---: | :---: | :---: |
| $\mathbf{1}$ | 0 | 0 | 0 | 10 | 0 | 10 |
| $\mathbf{0}$ | 1 | 0 | 10 | 0 | 0 | 0 |
| $\mathbf{0}$ | 0 | 1 | 8 | 0 | 10 | 8 |

unit vector for slack variable

## subdivide facet using slack variables

equilibrium iff all labels 1...m



## subdivide facet using slack variables

equilibrium iff all labels 1...m
4/35 $\quad 1 / 35 \quad 30 / 35$

| $\begin{array}{r} 4 / 5 \quad 1 / 5 \quad 0 \\ \text { (4) } 7 \text { (1) } \end{array}$ |  |  |  | 2 | 40/35 |
| :---: | :---: | :---: | :---: | :---: | :---: |
|  |  |  |  |  |  |
| (1) | 0 | 10 | 1 |  |  |
| (2) | 10 | 0 | 0 | 8 | 40/35 |
| (3) | 8 | 8 | 0 | 8 | 40/35 |



## The full dual construction

Given: nondegenerate $m \times n$ game $(A, B), m \leq n$.
Let $\mathrm{P}^{\Delta}=\operatorname{conv}[-\mathrm{M} \mathrm{Im}, \mathrm{B}]$, simplicial polytope.
Subdivide surface of $\mathrm{P}^{\Delta}$ into regions with labels $1, \ldots, \mathrm{~m}$ using columns of [ $\mathrm{Im}, \mathrm{A}$ ] that correspond to facets.

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Equilibrium ( $\mathrm{x}, \mathrm{y}$ )
$=$ point $f(y)$ on facet corresponding to $x$, vertices $=$ best responses player 2
with all labels $1, \ldots, \mathrm{~m}=$
best responses player 1

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Equilibrium ( $\mathrm{x}, \mathrm{y}$ )
$=$ point $f(y)$ on facet corresponding to $x$,
vertices $=$ best responses player 2 , unused strategies player 1
with all labels $1, \ldots, \mathrm{~m}=$
best responses player 1, slacks (=worse pay) for unused strategies.

## The full dual construction



Equilibria have all m labels


## Index = orientation



## Lemke-Howson paths



## Opposite index of endpoints



## Completely mixed NE of $3 \times 3$ coordination game made unique

$$
\left.\mathrm{A}=\left\lvert\, \begin{array}{llllll}
2 & 1 & 1 & 0 & 0 & 1 \\
1 & 2 & 1 & 1 & 0 & 0 \\
1 & 1 & 2 & 0 & 1 & 0
\end{array}\right.\right]
$$

$$
\begin{array}{rrr|rrr}
2 & 1 & 1 & 2.8 & 0.8 & -0.4 \\
1 & 2 & 1 & -0.4 & 2.8 & 0.8 \\
1 & 1 & 2 & 0.8 & -0.4 & 2.8
\end{array}
$$










## Adding new columns: polytope view





## Summary

## Dual construction

- Points = payoff columns (replies) of player 2
- Facet subdivision = player 1's best replies
- Visualization and characterization of index
- Illustration of L-H algorithm
- Low dimension m-1, for any $n$


## Other applications

- Components of equilibria, hyperstability
- Fixed point theorems, Sperner's lemma via bimatrix games


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